## - Writing an assembly language program

- Disassembly of MC9S12 op codes
- Use flow charts to lay out structure of program
- Use common flow structures
- if-then
- if-then-else
- do-while
- while
- Do not use spaghetti code
- Plan structure of data in memory
- Plan overall structure of program
- Work down to more detailed program structure
- Implement structure with instructions
- Optimize program to make use of instruction efficiencies
- Do not sacrifice clarity for efficiency

Binary, Hex and Decimal Numbers (4-bit representation)

| Binary | Hex | Decimal |
| :---: | :---: | :---: |
| 0000 | 0 | 0 |
| 0001 | 1 | 1 |
| 0010 | 2 | 2 |
| 0011 | 3 | 3 |
| 0100 | 4 | 4 |
| 0101 | 5 | 5 |
| 0110 | 6 | 6 |
| 0111 | 7 | 7 |
| 1000 | 8 | 8 |
| 1001 | 9 | 9 |
| 1010 | A | 10 |
| 1011 | B | 11 |
| 1100 | C | 12 |
| 1101 | D | 13 |
| 1110 | E | 14 |
| 1111 | F | 15 |

What does a number represent?
Binary numbers are a code, and represent what the programmer intends for the code.
0x72 Some possible meanings:
'r' (ASCII)
INC MEM (hh ll) (HC12 instruction)
2.26 V (Input from A/D converter)
$114_{10}$ (Unsigned number)
$+114_{10}$ (Signed number)
Set temperature in room to $69^{\circ} \mathrm{F}$
Set cruise control speed to 120 mph

Binary to Unsigned Decimal:
Convert Binary to Unsigned Decimal
$1111011_{2}$
$1 \times 2^{6}+1 \times 2^{5}+1 \times 2^{4}+1 \times 2^{3}+0 \times 2^{2}+1 \times 2^{1}+1 \times 2^{0}$
$1 \times 64+1 \times 32+1 \times 16+1 \times 8+0 \times 4+1 \times 2+1 \times 1$
12310

Hex to Unsigned Decimal
Convert Hex to Unsigned Decimal
82D6 16
$8 \times 16^{3}+2 \times 16^{2}+13 \times 16^{1}+6 \times 16^{0}$
$8 \times 4096+2 \times 256+13 \times 16+6 \times 1$
$33494{ }_{10}$

Unsigned Decimal to Hex
Convert Unsigned Decimal to Hex

| Division | Q | $\mathbf{R}$ |  |
| :---: | :---: | :---: | :---: |
|  |  | Decimal | Hex |
| $721 / 16$ | 45 | 1 | 1 |
| $45 / 16$ | 2 | 13 | D |
| $2 / 16$ | 0 | 2 | 2 |

$721_{10}=2$ D1 $_{16}$

Signed Number Representation in 2's Complement Form:

If the most significant bit (MSB) is 0 (most significant hex digit $0-7$ ), then the number is positive.
Get decimal equivalent by converting number to decimal, and use the + sign.

## Example for 8-bit number:

$$
\begin{aligned}
\mathbf{3} \mathbf{A}_{16}-> & +\left(3 \times 16^{1}+10 \times 16^{0}\right)_{10} \\
& +(3 \times 16+10 \times 1)_{10} \\
& +\mathbf{5 8} \mathbf{1 0}
\end{aligned}
$$

If the most significant bit is 1 (most significant hex digit $8-\mathrm{F}$ ), then the number is negative.
Get decimal equivalent by taking 2's complement of number, converting to decimal, and using - sign.

Example for 8-bit number:
$\mathbf{A 3}_{16}$-> - (5D) ${ }_{16}$
$-\left(5 \times 16^{1}+13 \times 16^{0}\right)_{10}$

- $(5 \times 16+13 \times 1)_{10}$
- 9310

One's complement table makes it simple to finding 2's complements
One's Complement Table

| 0 | F |
| :---: | :---: |
| 1 | E |
| 2 | D |
| 3 | C |
| 4 | B |
| 5 | A |
| 6 | 9 |
| 7 | 8 |
| 7 |  |

To take two's complement, add one to one's complement.
Take two's complement of D0C3:

$$
2 \mathrm{~F} 3 \mathrm{C}+1=\mathbf{2 F} \mathbf{D D}
$$

Addition and Subtraction of Binary and Hexadecimal Numbers
Setting the C (Carry), V (Overflow), N (Negative) and Z (Zero) bits

How the $\mathbf{C ,}, \mathbf{V}, \mathbf{N}$ and Z bits of the CCR are changed
N bit is set if result of operation is negative $(\mathrm{MSB}=1)$
Z bit is set if result of operation is zero $($ All bits $=0)$
V bit is set if operation produced an overflow

C bit is set if operation produced a carry (borrow on subtraction)

Note: Not all instructions change these bits of the CCR

## Addition of Hexadecimal Numbers

## ADDITION:

C bit set when result does not fit in word

V bit set when $\mathrm{P}+\mathrm{P}=\mathrm{N}$ or $\mathrm{N}+\mathrm{N}=\mathrm{P}$
N bit set when MSB of result is 1

Z bit set when result is 0

| 7A | 2A | AC | AC |
| :---: | :---: | :---: | :---: |
| +52 | +52 | +8A | +72 |
| CC | 7 C | 36 | 1E |
| C: 0 | C: 0 | C: 1 | C: 1 |
| $\mathrm{V}: 1$ | V: 0 | $\mathrm{V}: 1$ | V: 0 |
| $\mathrm{N}: 1$ | N: 0 | $\mathrm{N}: 0$ | $\mathrm{N}: 0$ |
| Z: 0 | Z: 0 | Z: 0 | Z: 0 |

## Subtraction of Hexadecimal Numbers

## SUBTRACTION:

C bit set on borrow (when the magnitude of the subtrahend is greater than the minuend

V bit set when $\mathrm{N}-\mathrm{P}=\mathrm{P}$ or $\mathrm{P}-\mathrm{N}=\mathrm{N}$

N bit set when MSB is 1
Z bit set when result is 0

| 7 A | 2 A | AC | AC |
| ---: | ---: | ---: | :---: |
| -5 C | -5 C | -8 A | -72 |
| ----------- | --- | 32 |  |


| C: 0 | C: | C: 0 | C: 0 |
| :--- | :--- | :--- | :--- |
| V: 0 | V: 0 | V: 0 | V: 1 |
| N: 0 | N: 1 | N: 0 | N: 0 |
| Z: 0 | Z: 0 | Z: 0 | Z: 0 |

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Writing Assembly Language Programs
Use Flowcharts to Help Plan Program Structure

## Flow chart symbols:



## IF-THEN Flow Structure



## if (C)

\{
A;
\}

EXAMPLE:

| if ( $\mathrm{C}<10$ ) |  | CMPA | \#10 ; if (A<10) |
| :---: | :---: | :---: | :---: |
| 1 |  | BLT | L1 ; signed numbers |
| var $=5 ;$ |  | BRA | L2 |
| \} | L1: | LDAB | \#5 ; var=5 |
|  |  | STAB | var |
|  |  | next ins |  |

OR:

| CMPA | \#10 | ; if $(\mathrm{A}<10)$ |
| :--- | :--- | :--- |
| BGE | L2 | $;$ signed numbers |
| LDAB | \#5 | ; var=5 |
| STAB | var |  |
| next instruction |  |  |

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## IF-THEN-ELSE Flow Structure


if(A < 10)
\{ $\operatorname{var}=5 ;$
\}
else
\{
$\operatorname{var}=0 ;$
\}

|  | CMPA | $\# 10 ;$ if $(\mathrm{A}<10)$ |
| :--- | :--- | :--- |
|  | BLT | L1 $;$ signed numbers |
|  | CLR | var ; var=0 |
|  | BRA | L2 |
| L1: | LDAB | $\# 5 \quad$; var=5 |
|  | STAB | var |
| L2: | next instruction |  |

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DO WHILE Flow Structure


EXAMPLE:

| $\mathrm{i}=0$; |  | LDX | \#table |  |
| :---: | :---: | :---: | :---: | :---: |
| do |  | CLRA |  | ; i=0 |
| , | L1: | ASR | 1,X+ | ; table[i] /=2 |
| table[i]=table[i]/2; |  | INCA |  | ; i=i+1 |
| $i=i+1$ |  | CMPA | \#LEN | ; while(i<=10) |
| \} |  | BLE | L1 | ; unsigned |
| while (i <= LEN); |  |  |  | ; numbers |

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WHILE Flow Structure

while (C)
\{
A;
\}

## EXAMPLE:

| $\begin{aligned} & \mathbf{i}=\mathbf{0} \\ & \text { while( } \mathrm{i}<=\text { LEN }) \end{aligned}$ |
| :---: |
| $\left\{\begin{array}{l} \text { table[i]=table[i]*2; } \\ \substack{\mathrm{i}=\mathrm{i}+1 ;} \end{array}\right.$ |
| ) |


|  | LDX | \#table |
| :--- | :--- | :--- |
|  | CLRA |  |
| L1: | CMPA | \#LEN |
|  | BLT | L2 |
|  | BRA | L3 |
| L2: | ASL | 1, X+ |
|  | INCA |  |
|  | BRA | L1 |
| L3: | next instruction |  |

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Use Good Structure When Writing Programs

- Do Not Use Spaghetti Code



## Example Program: Divide a table of data by 2

Problem: Start with a table of data. The table consists of 5 values. Each value is between 0 and 255. Create a new table whose contents are the original table divided by 2.

1. Determine where code and data will go in memory.

Code at $\$ 2000$, data at $\$ 1000$.
2. Determine type of variables to use.

Because data will be between 0 and 255, can use unsigned 8 -bit numbers.
3. Draw a picture of the data structures in memory:

4. Strategy: Because we are using a table of data, we will need pointers to each table so we can keep track of which table element we are working on.

Use the $X$ and $Y$ registers as pointers to the tables.

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5. Use a simple flow chart to plan structure of program.

6. Need a way to determine when we reach the end of the table.

One way: Use a counter (say, register A) to keep track of how many Elements we have processed.

7. Add code to implement blocks:

8. Write the program:
; Program to divide a table by two
; and store the results in memory

| prog: equ <br> data: equ | $\begin{aligned} & \$ 2000 \\ & \$ 1000 \end{aligned}$ |  |
| :---: | :---: | :---: |
| count: equ | 5 |  |
| org | prog | ; Set program counter to 0x2000 |
| Idaa | \#count | ; Use A as counter |
| ldx | \#table1 | ; Use X as data pointer to table1 |
| ldy | \#table2 | ; Use Y as data pointer to table2 |
| 11: ldab | 0,x | ; Get entry from table1 |
| lsrb |  | ; Divide by two (unsigned) |
| stab | 0,y | ; Save in table2 |
| inx |  | ; Increment table1 pointer |
| iny |  | ; Increment table 2 pointer |
| deca |  | ; Decrement counter |
| bne | 11 | ; Counter ! $=0=>$ more entries to divide |
| swi |  | ; Done |

org data
table1: dc.b \$07,\$c2,\$3a,\$68,\$f3
table2: ds.b count
9. Advanced: Optimize program to make use of instructions set efficiencies:

```
; Program to divide a table by two
; and store the results in memory
prog: equ 
count: equ 5
\begin{tabular}{llll} 
& org & prog & ; Set program counter to 0x1000 \\
& ldaa & \#count & ; Use A as counter \\
& ldx & \#table1 & ; Use X as data pointer to table1 \\
l1: & ldy & \#table2 & ; Use Y as data pointer to table2 \\
ldab & \(\mathbf{1 , x +}\) & ; Get entry from table1; then inc pointer \\
& lsrb & & ; Divide by two (unsigned) \\
& stab & \(\mathbf{1 , y +}\) & ; Save in table2; then inc pointer \\
& dbne & \(\mathbf{a , 1 1}\) & ; Decrement counter; if not 0, more to do \\
& swi & & ; Done
\end{tabular}
org data
table1: dc.b \$07,\$c2,\$3a,\$68,\$f3
table2: ds.b count
```

EE 308 Spring 2011

## TOP-DOWN PROGRAM DESIGN

- PLAN DATA STRUCTURES IN MEMORY
- START WITH A LARGE PICTURE OF THE PROGRAM STRUCTURE
- WORK DOWN TO MORE DETAILED STRUCTURE
- TRANSLATE STRUCTURE INTO CODE
- OPTIMIZE FOR EFFICIENCY

DO NOT SACRIFICE CLARITY FOR EFFICIENCY

