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- Decimal, Hexadecimal and Binary Numbers
- Writing an assembly language program
- Disassembly of MC9S12 op codes
- Use flow charts to lay out structure of program
o Use common flow structures
- if-then
- if-then-else
- do-while
- while
- Do not use spaghetti code
- Plan structure of data in memory
- Plan overall structure of program
- Work down to more detailed program structure
- Implement structure with instructions
- Optimize program to make use of instruction efficiencies
- Do not sacrifice clarity for efficiency

Binary, Hex and Decimal Numbers (4-bit representation)

| Binary | Hex | Decimal |
| :---: | :---: | :---: |
| 0000 | 0 | 0 |
| 0001 | 1 | 1 |
| 0010 | 2 | 2 |
| $\ldots$ | $\ldots$ | $\ldots$ |
| 1010 | A | 10 |
| 1011 | B | 11 |
| 1100 | C | 12 |
| 1101 | D | 13 |
| 1110 | E | 14 |
| 1111 | F | 15 |

## What does a number represent?

Binary numbers are a code, and represent what the programmer intends for the code.

0x72 Some possible meanings:
'r' (ASCII)
INC MEM (hh ll) (HC12 instruction)
2.26V (Input from A/D converter)
$114_{10}$ (Unsigned number)
$+114_{10}$ (Signed number)
Set temperature in room to $69^{\circ} \mathrm{F}$
Set cruise control speed to 120 mph

## Binary to Unsigned Decimal:

Convert Binary to Unsigned Decimal $1111011_{2}$
$1 \times 2^{6}+1 \times 2^{5}+1 \times 2^{4}+1 \times 2^{3}+0 \times 2^{2}+1 \times 2^{1}+1 \times 2^{0}$
$1 \times 64+1 \times 32+1 \times 16+1 \times 8+0 \times 4+1 \times 2+1 \mathrm{x} 1$
$123{ }_{10}$

Hex to Unsigned Decimal
Convert Hex to Unsigned Decimal
82D6 ${ }_{16}$
$8 \times 16^{3}+2 \times 16^{2}+13 \times 16^{1}+6 \times 16^{0}$
$8 \times 4096+2 \times 256+13 \times 16+6 \times 1$
$33494{ }_{10}$

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## Unsigned Decimal to Hex

Convert Unsigned Decimal to Hex

| Division | Q | $\mathbf{R}$ |  |
| :---: | :---: | :---: | :---: |
|  |  | Decimal | Hex |
| $721 / 16$ | 45 | 1 | 1 |
| $45 / 16$ | 2 | 13 | D |
| $2 / 16$ | 0 | 2 | 2 |

$721_{10}=2 D 1_{16}$

Signed Number Representation in 2's Complement Form:

If the most significant bit (MSB) is 0 (most significant hex digit $0-7$ ), then the number is positive.

Get decimal equivalent by converting number to decimal, and use the + sign.

## Example for 8-bit number:

$$
\begin{aligned}
\mathbf{3} \mathbf{A}_{16}-> & +\left(3 \times 16^{1}+10 \times 16^{0}\right)_{10} \\
& +(3 \times 16+10 \times 1)_{10} \\
& +58 \mathbf{1 0}
\end{aligned}
$$

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If the most significant bit is 1 (most significant hex digit $8-\mathrm{F}$ ), then the number is negative.

Get decimal equivalent by taking 2's complement of number, converting to decimal, and using - sign.

Example for 8-bit number:

$$
\begin{aligned}
\mathbf{A} \mathbf{3}_{16}- & -(5 C+1)_{16} \\
& -\left(5 \times 16^{1}+13 \times 16^{0}\right)_{10} \\
& -(5 \times 16+13 \times 1)_{10} \\
& -\mathbf{9 3} \mathbf{1 0}
\end{aligned}
$$

One's complement table makes it simple to finding 2's complements


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One's complement table makes it simple to finding 2's complements


To take two's complement, add one to one's complement.

Take two's complement of D0C3:

$$
2 \mathrm{~F} 3 \mathrm{C}+1=2 \mathrm{~F} 3 \mathrm{D}
$$

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# Addition and Subtraction of Binary and Hexadecimal Numbers 

Setting the C (Carry), V (Overflow), N (Negative) and Z (Zero) bits

How the $\mathrm{C}, \mathrm{V}, \mathrm{N}$ and Z bits of the CCR are changed?
N bit is set if result of operation is negative ( $\mathrm{MSB}=1$ )
Z bit is set if result of operation is zero (All bits $=0$ )
V bit is set if operation produced an overflow

C bit is set if operation produced a carry (borrow on subtraction)

Note: Not all instructions change these bits of the CCR

## Addition of Hexadecimal Numbers

## ADDITION:

C bit set when result does not fit in word
V bit set when $\mathrm{P}+\mathrm{P}=\mathrm{N}$ or

$$
\mathrm{N}+\mathrm{N}=\mathrm{P}
$$

N bit set when MSB of result is 1

Z bit set when result is 0

| 7A | 2A | AC | AC |
| :---: | :---: | :---: | :---: |
| +52 | +52 | +8A | +72 |
| CC | 7C | 36 | 1 E |
| C: 0 | C: 0 | C: 1 | C: 1 |
| V: 1 | V: 0 | $\mathrm{V}: 1$ | V: 0 |
| N: 1 | N: 0 | N: 0 | $\mathrm{N}: 0$ |
| Z: 0 | Z: 0 | Z: 0 | Z: 0 |

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## Subtraction of Hexadecimal Numbers

## SUBTRACTION:

C bit set on borrow (when the magnitude of the subtrahend is greater than the minuend

V bit set when $\mathrm{N}-\mathrm{P}=\mathrm{P}$ or

$$
\mathrm{P}-\mathrm{N}=\mathrm{N}
$$

N bit set when MSB is 1

Z bit set when result is 0

| 7A | 8A | 5C | 2C |
| :---: | :---: | :---: | :---: |
| -5C | -5C | -8A | -72 |
| ---- | 2E--- | D2 | BA |
| C: 0 | C: 0 | C: 1 | C: 1 |
| V: 0 | V: 1 | V: 1 | V: 0 |
| N: 0 | N: 0 | N: 1 | $\mathrm{N}: 1$ |
| Z: 0 | Z: 0 | Z: 0 | Z: 0 |

## Writing Assembly Language Programs

## Use Flowcharts to Help Plan Program Structure

## Flow chart symbols:



## IF-THEN Flow Structure



EXAMPLE:

```
if (A<10)
{
    var = 5;
    }
```

L1: LDAB \#5 ; var=5
STAB var
L2: next instruction

OR:
CMPA \#10; if(A<10)
BGE L2 ; signed numbers

LDAB \#5 ; var=5
STAB var
L2: next instruction

## IF-THEN-ELSE Flow Structure


if(A< 10)
if(A< 10)
{
{
var = 5;
var = 5;
}
}
else L1:
{
{
var = 0;
var = 0;
}
}

DO WHILE Flow Structure


EXAMPLE:

```
i=0;
do
{
    table[i]=table[i]/2;
    i=i+1;
}
while (i <= LEN);
LDX \#table
CLRA \(\quad\) i=0
L1: ASR 1,X+ ; table[i]/=2
INCA \(\quad ; \mathrm{i}=\mathrm{i}+1\)
CMPA \#LEN ; while(i<=10)
BLE L1 ; unsigned
; numbers
```


## WHILE Flow Structure



## EXAMPLE:

| $\begin{aligned} & \mathrm{i}=\mathbf{0} \text {; } \\ & \text { while( } \mathrm{i} \text { <= LEN) } \end{aligned}$ |  | LDX CLRA | \#table |
| :---: | :---: | :---: | :---: |
| \{ | L1: | CMPA | \#LEN |
| table[i]=table[i]*2; |  | BLT | L2 |
| $\mathrm{i}=\mathbf{i}+1$; |  | BRA | L3 |
| \} | L2: | ASL <br> INCA | 1,X+ |
|  |  | BRA | L1 |
|  | L3: | next i | uction |

## Use Good Structure When Writing Programs

## - Do Not Use Spaghetti Code

## SPAGHEITI CUE

DO NOT USE


## Example Program: Divide a table of data by 2

Problem: Start with a table of data. The table consists of 5 values. Each value is between 0 and 255 . Create a new table whose contents are the original table divided by 2.

1. Determine where code and data will go in memory. Code at $\$ 2000$, data at $\$ 1000$.
2. Determine type of variables to use.

Because data will be between 0 and 255, can use unsigned 8bit numbers.
3. Draw a picture of the data structures in memory:

```
table1:
```

table2:


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4. Strategy: Because we are using a table of data, we will need pointers to each table so we can keep track of which table element we are working on.

## Use the $X$ and $Y$ registers as pointers to the tables.

5. Use a simple flow chart to plan structure of program.

6. Need a way to determine when we reach the end of the table.

One way: Use a counter (say, register A) to keep track of how many Elements we have processed.


## 7. Add code to implement blocks:


8. Write the program:
; Program to divide a table by two
; and store the results in memory
prog: equ $\$ 2000$
data: equ $\$ 1000$
count: equ 5
org prog ; Set program counter to 0x2000
Idaa \#count ; Use A as counter
Idx \#table1 ; Use $X$ as data pointer to table1
Idy \#table2 ; Use Y as data pointer to table2
11:

| Idab | $\mathbf{0 , x}$ | ; Get entry from table1 |
| :--- | :--- | :--- |
| lsrb |  | ; Divide by two (unsigned) |
| stab | $\mathbf{0 , y}$ | ; Save in table2 |
| inx |  | ; Increment table1 pointer |
| iny |  | Increment table2 pointer |
| deca | I Decrement counter |  |
| bne | $\mathbf{1 1}$ | ; Counter ! $=0=>$ more entries |
|  |  | to divide |
| swi | Done |  |

org data
table1: dc.b \$07,\$c2,\$3a,\$68,\$f3
table2: ds.b count

## 9. Advanced: Optimize program to make use of instructions set

 efficiencies:; Program to divide a table by two
; and store the results in memory

| prog: equ data: | $\begin{aligned} & \$ 100 \\ & \text { equ } \end{aligned}$ | \$2000 |  |
| :---: | :---: | :---: | :---: |
| count: | equ | 5 |  |
|  | org | prog | ; Set program counter to 0x1000 |
|  | ldaa | \#count | ; Use A as counter |
|  | ldx | \#table1 | ; Use $X$ as data pointer to table1 |
|  | ldy | \#table2 | ; Use Y as data pointer to table2 |
| 11: | ldab <br> lsrb | 1,x+ | ; Get entry from table1; then inc ptr. <br> ; Divide by two (unsigned) |
|  | stab | 1, ${ }^{+}$ | ; Save in table2; then inc ptr. |
|  | dbne | a,11 | ; Decrement counter; if not 0 , <br> ; more to do |
|  | swi |  | ; Done |


|  | org | data |
| :--- | :--- | :--- |
| table1: | dc.b | $\$ 07, \$ c 2, \$ 3 a, \$ 68, \$ f 3$ |

table2: ds.b count

- PLAN DATA STRUCTURES IN MEMORY
- START WITH A LARGE PICTURE OF THE PROGRAM STRUCTURE
- WORK DOWN TO MORE DETAILED STRUCTURE
- TRANSLATE STRUCTURE INTO CODE
- OPTIMIZE FOR EFFICIENCY

DO NOT SACRIFICE CLARITY FOR EFFICIENCY

