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- Decimal, Hexadecimal and Binary Numbers
- Writing an assembly language program
- Disassembly of MC9S12 op codes
- Use flow charts to lay out structure of program
o Use common flow structures
- if-then
- if-then-else
- do-while
- while
- Do not use spaghetti code
- Plan structure of data in memory
- Plan overall structure of program
- Work down to more detailed program structure
- Implement structure with instructions
- Optimize program to make use of instruction efficiencies
- Do not sacrifice clarity for efficiency

Binary, Hex and Decimal Numbers (4-bit representation)

| Binary | Hex | Decimal |
| :---: | :---: | :---: |
| 0000 | 0 | 0 |
| 0001 | 1 | 1 |
| 0010 | 2 | 2 |
| $\ldots$ | $\ldots$ | $\ldots$ |
| 1010 | A | 10 |
| 1011 | B | 11 |
| 1100 | C | 12 |
| 1101 | D | 13 |
| 1110 | E | 14 |
| 1111 | F | 15 |

## What does a number represent?

Binary numbers are a code, and represent what the programmer intends for the code.

0x72 Some possible meanings:
'r' (ASCII)
INC MEM (hh ll) (HC12 instruction)
2.26 V (Input from A/D converter)
$114_{10}$ (Unsigned number)
$+114_{10}$ (Signed number)
Set temperature in room to $69^{\circ} \mathrm{F}$
Set cruise control speed to 120 mph

## Binary to Unsigned Decimal:

Convert Binary to Unsigned Decimal $1111011_{2}$
$1 \times 2^{6}+1 \times 2^{5}+1 \times 2^{4}+1 \times 2^{3}+0 \times 2^{2}+1 \times 2^{1}+1 \times 2^{0}$
$1 \times 64+1 \times 32+1 \times 16+1 \times 8+0 \times 4+1 \times 2+1 \mathrm{x} 1$
$123{ }_{10}$

Hex to Unsigned Decimal
Convert Hex to Unsigned Decimal
82D6 ${ }_{16}$
$8 \times 16^{3}+2 \times 16^{2}+13 \times 16^{1}+6 \times 16^{0}$
$8 \times 4096+2 \times 256+13 \times 16+6 \times 1$
$33494{ }_{10}$

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## Unsigned Decimal to Hex

Convert Unsigned Decimal to Hex

| Division | Q | R |  |
| :---: | :---: | :---: | :---: |
|  |  | Decimal | Hex |
| $721 / 16$ | 45 | 1 | 1 |
| $45 / 16$ | 2 | 13 | $D$ |
| $2 / 16$ | 0 | 2 | 2 |

$$
721_{10}=2 D 1_{16}
$$

Signed Number Representation in 2's Complement Form:

If the most significant bit (MSB) is 0 (most significant hex digit $0-7$ ), then the number is positive.

Get decimal equivalent by converting number to decimal, and use the + sign.

## Example for 8-bit number:

$$
\begin{aligned}
\mathbf{3} \mathbf{A}_{16}-> & +\left(3 \times 16^{1}+10 \times 16^{0}\right)_{10} \\
& +(3 \times 16+10 \times 1)_{10} \\
& +58 \mathbf{x p}_{10}
\end{aligned}
$$

If the most significant bit is 1 (most significant hex digit $8-F$ ), then the number is negative.

Get decimal equivalent by taking 2's complement of number, converting to decimal, and using - sign.

Example for 8-bit number:

$$
\begin{aligned}
\mathbf{A} \mathbf{3}_{16}- & -(5 C+1)_{16} \\
& -\left(5 \times 16^{1}+13 \times 16^{0}\right)_{10} \\
& -(5 \times 16+13 \times 1)_{10} \\
& -\mathbf{9 3} \mathbf{1 0}
\end{aligned}
$$

One's complement table makes it simple to finding 2's complements


One's complement table makes it simple to finding 2's complements


To take two's complement, add one to one's complement.
Take two's complement of D0C3:

$$
2 \mathrm{~F} 3 \mathrm{C}+1=2 \mathrm{~F} 3 \mathrm{D}
$$

# Addition and Subtraction of Binary and Hexadecimal Numbers 

Setting the C (Carry), V (Overflow), N (Negative) and Z (Zero) bits

How the $\mathrm{C}, \mathrm{V}, \mathrm{N}$ and Z bits of the CCR are changed?
N bit is set if result of operation is negative $(\mathrm{MSB}=1)$
Z bit is set if result of operation is zero (All bits $=0$ )

V bit is set if operation produced an overflow

C bit is set if operation produced a carry (borrow on subtraction)

Note: Not all instructions change these bits of the CCR

## Addition of Hexadecimal Numbers

## ADDITION:

C bit set when result does not fit in word
V bit set when $\mathrm{P}+\mathrm{P}=\mathrm{N}$ or

$$
\mathrm{N}+\mathrm{N}=\mathrm{P}
$$

N bit set when MSB of result is 1

Z bit set when result is 0

| 7A | 2A | AC | AC |
| :---: | :---: | :---: | :---: |
| +52 | +52 | +8A | +72 |
| CC | 7C | 36 | 1E |
| C: 0 | C: 0 | C: 1 | C: 1 |
| V: 1 | V: 0 | V: 1 | V: 0 |
| $\mathrm{N}: 1$ | N: 0 | N: 0 | N: 0 |
| Z: 0 | Z: 0 | Z: 0 | Z: 0 |

## Subtraction of Hexadecimal Numbers

## SUBTRACTION:

C bit set on borrow (when the magnitude of the subtrahend is greater than the minuend

V bit set when $\mathrm{N}-\mathrm{P}=\mathrm{P}$ or

$$
\mathrm{P}-\mathrm{N}=\mathrm{N}
$$

N bit set when MSB is 1

Z bit set when result is 0

| 7A | 8A | 5C | 2C |
| :---: | :---: | :---: | :---: |
| -5C | -5C | -8A | -72 |
| ----- | 2E | ------ | BA |
| C: 0 | C: 0 | C: 1 | C: 1 |
| V: 0 | V: 1 | V: 1 | V: 0 |
| N: 0 | N: 0 | $\mathrm{N}: 1$ | $\mathrm{N}: 1$ |
| Z: 0 | Z: 0 | Z: 0 | Z: 0 |

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## Writing Assembly Language Programs

## Use Flowcharts to Help Plan Program Structure

## Flow chart symbols:



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## IF-THEN Flow Structure



## if (C) <br> \{ <br> A; <br> \}

EXAMPLE:

```
if (A<10)
{
    var = 5;
    }
```

    CMPA \#10 ; if ( \(\mathrm{A}<10\) )
    BLT L1 ; signed numbers
    BRA L2
    L1: LDAB \#5 ; var=5
STAB var
L2: next instruction
OR:
CMPA \#10; if(A<10)
BGE L2 ; signed numbers
LDAB \#5 ; var=5
STAB var
L2: next instruction

## IF-THEN-ELSE Flow Structure



DO WHILE Flow Structure


EXAMPLE:

```
i=0;
do
{
    table[i]=table[i]/2;
    i=i+1;
}
while (i <= LEN);
LDX \#table
CLRA ;i=0
L1: ASR 1,X+ ; table[i]/=2
INCA ; i=i+1
CMPA \#LEN ; while(i<=10)
BLE L1 ; unsigned
; numbers
```


## WHILE Flow Structure



EXAMPLE:

| $\mathbf{i}=0 ;$ |  | LDX | \#table |
| :---: | :---: | :---: | :---: |
| while( i <= LEN) |  | CLRA |  |
| \{ | L1: | CMPA | \#LEN |
| table[i]=table[i]*2; |  | BLT | L2 |
| $\mathrm{i}=\mathbf{i}+1$; |  | BRA | L3 |
| \} | L2: | ASL <br> INCA | 1,X+ |
|  |  | BRA | L1 |
|  | L3: | next i | uction |

## Use Good Structure When Writing Programs

## - Do Not Use Spaghetti Code

SPAGHEITI CUE
DO NOT USE


## Example Program: Divide a table of data by 2

Problem: Start with a table of data. The table consists of 5 values. Each value is between 0 and 255 . Create a new table whose contents are the original table divided by 2.

1. Determine where code and data will go in memory. Code at $\$ 2000$, data at $\$ 1000$.
2. Determine type of variables to use.

Because data will be between 0 and 255, can use unsigned 8bit numbers.
3. Draw a picture of the data structures in memory:

4. Strategy: Because we are using a table of data, we will need pointers to each table so we can keep track of which table element we are working on.

## Use the $X$ and $Y$ registers as pointers to the tables.

5. Use a simple flow chart to plan structure of program.

6. Need a way to determine when we reach the end of the table.

One way: Use a counter (say, register A) to keep track of how many Elements we have processed.

7. Add code to implement blocks:

8. Write the program:
; Program to divide a table by two
; and store the results in memory
prog: equ $\$ 2000$
data: equ $\$ 1000$
count: equ 5
org prog ; Set program counter to $0 \times 2000$
Idaa \#count ; Use A as counter
ldx \#table1 ; Use $X$ as data pointer to table1
ldy \#table2 ; Use Y as data pointer to table2
11:
ldab 0,x ; Get entry from table1
lsrb ; Divide by two (unsigned)
stab 0,y ; Save in table2
inx ; Increment table1 pointer
iny ; Increment table2 pointer
deca ; Decrement counter
bne $\mathbf{1 1}$; Counter ! $=0=>$ more entries
; to divide
swi ; Done
org data
table1: dc.b \$07,\$c2,\$3a,\$68,\$43
table2: ds.b count
9. Advanced: Optimize program to make use of instructions set efficiencies:

```
; Program to divide a table by two
; and store the results in memory
\begin{tabular}{lll} 
prog: & equ & \(\$ 1000\) \\
data: & equ & \(\$ 2000\)
\end{tabular}
count: equ 5
org prog ; Set program counter to 0x1000
ldaa #count ; Use A as counter
ldx #table1 ; Use X as data pointer to table1
ldy #table2 ; Use Y as data pointer to table2
11: Idab 1,x+ ; Get entry from table1; then inc ptr.
    lsrb ; Divide by two (unsigned)
stab 1,\mp@subsup{\mathbf{y}}{}{+}\quad; Save in table2; then inc ptr.
dbne a,l1 ; Decrement counter; if not 0,
            ; more to do
swi ; Done
```

|  | org | data |
| :--- | :--- | :--- |
| table1: | dc.b | $\$ 07, \$ c 2, \$ 3 a, \$ 68, \$ f 3$ |

table2: ds.b count

## TOP-DOWN PROGRAM DESIGN

- PLAN DATA STRUCTURES IN MEMORY
- START WITH A LARGE PICTURE OF THE PROGRAM STRUCTURE
- WORK DOWN TO MORE DETAILED STRUCTURE
- TRANSLATE STRUCTURE INTO CODE
- OPTIMIZE FOR EFFICIENCY

DO NOT SACRIFICE CLARITY FOR EFFICIENCY

