

- **MC9S12 Assembler Directives**
- **A Summary of MC9S12 Instructions**
- **Disassembly of MC9S12 op codes**
 - Number of cycles and time taken to execute an MC9S12 program
 - A label is a name assigned to the address of the location counter where the label is defined
 - Use of assembler directives
 - A summary of MC9S12 instruction
 - How to disassemble an MC9S12 instruction sequence

Summary of HCS12 addressing modes

ADDRESSING MODES

Name	Example	Op Code	Effective Address
INH Inherent	ABA	18 06	None
IMM Immediate	LDAA #\$35	86 35	PC + 1
DIR Direct	LDAA \$35	96 35	0x0035
EXT Extended	LDAA \$2035	B6 20 35	0x2035
IDX Indexed	LDAA 3, X	A6 03	X + 3
IDX1	LDAA 30, X	A6 E0 13	X + 30
IDX2	LDAA 300, X	A6 E2 01 2C	X + 300
IDX Indexed Postincrement	LDAA 3, X+	A6 32	X (X+3 -> X)
IDX Indexed Preincrement	LDAA 3, +X	A6 22	X+3 (X+3 -> X)
IDX Indexed Postdecrement	LDAA 3, X-	A6 3D	X (X-3 -> X)
IDX Indexed Predecrement	LDAA 3, -X	A6 2D	X-3 (X-3 -> X)
REL Relative	BRA \$1050 LBRA \$1F00	20 23 18 20 0E CF	PC + 2 + Offset PC + 4 + Offset

A few instructions have two effective addresses:

- **MOVB # $\$AA$, $\$1C00$** Move byte 0xAA (IMM) to address \$1C00 (EXT)
- **MOVW 0,X,0,Y** Move word from address pointed to by X (IDX) to address pointed to by Y (IDX)

A few instructions have three effective addresses:

- **BRSET FOO,# $\$03$,LABEL** Branch to LABEL (REL) if bits # $\$03$ (IMM) of variable FOO (EXT) are set.

Using X and Y as Pointers

- Registers X and Y are often used to point to data.
- To initialize pointer use

ldx #table
not

ldx table

- For example, the following loads the address of table (\$1000) into X; i.e., X will point to table:

ldx #table ; Address of table \Rightarrow X

The following puts the first two bytes of table (\$0C7A) into X. X will **not** point to table:

ldx table ; First two bytes of table \Rightarrow X

- To step through table, need to increment pointer after use

ldaa 0,x
inx

or

ldaa 1,x+

table	0C
	7A
	D5
	00
	61
	62
	63
	64

```
table:  org    $900
        dc.b  12,122,-43,0
        dc.b  'a','b','c','d'
```

Which branch instruction should you use?

Branch if $A > B$

Is $0xFF > 0x00$?

If unsigned, $0xFF = 255$ and $0x00 = 0$,
so $0xFF > 0x00$

If signed, $0xFF = -1$ and $0x00 = 0$,
so $0xFF < 0x00$

Using unsigned numbers: **BHI** (checks C bit of CCR)

Using signed numbers: **BGT** (checks V bit of CCR)

For unsigned numbers, use branch instructions which check C bit

For signed numbers, use branch instructions which check V bit

Hand Assembling a Program

To hand-assemble a program, do the following:

1. Start with the **org** statement, which shows where the first byte of the program will go into memory.

(e.g., **org \$2000** will put the first instruction at address **\$2000**.)

2. Look at the first instruction. Determine the addressing mode used.

(e.g., **ldab #10** uses IMM mode.)

3. Look up the instruction in the **MC9S12 S12CPUV2 Reference Manual**, find the appropriate Addressing Mode, and the Object Code for that addressing mode.

(e.g., **ldab IMM** has object code **C6 ii**.)

- **Table A.1 of the S12CPUV2 Reference Manual** has a concise summary of the instructions, addressing modes, op-codes, and cycles.

4. Put in the object code for the instruction, and put in the appropriate operand. Be careful to convert decimal operands to hex operands if necessary.

(e.g., **ldab #10** becomes **C6 0A**.)

5. Add the number of bytes of this instruction to the address of the instruction to determine the address of the next instruction.

(e.g., **\$2000 + 2 = \$2002** will be the starting address of the next instruction.)

```
org $2000
ldab #10
loop: clra
      dbne b,loop
      swi
```

Table A-1. Instruction Set Summary (Sheet 7 of 14)

Source Form	Operation	Addr. Mode	Machine Coding (hex)	Access Detail		S X H I	N Z V C
				HCS12	M68HC12		
LBGIT <i>rel16</i>	Long Branch if Greater Than (if Z + (N ⊕ V) = 0) (signed)	REL	18 2E qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBHI <i>rel16</i>	Long Branch if Higher (if C + Z = 0) (unsigned)	REL	18 22 qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBHS <i>rel16</i>	Long Branch if Higher or Same (if C = 0) (unsigned) same function as LBCC	REL	18 24 qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBLE <i>rel16</i>	Long Branch if Less Than or Equal (if Z + (N ⊕ V) = 1) (signed)	REL	18 2F qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBLO <i>rel16</i>	Long Branch if Lower (if C = 1) (unsigned) same function as LBCCS	REL	18 25 qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBSL <i>rel16</i>	Long Branch if Lower or Same (if C + Z = 1) (unsigned)	REL	18 23 qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBLT <i>rel16</i>	Long Branch if Less Than (if N ⊕ V = 1) (signed)	REL	18 2D qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBMI <i>rel16</i>	Long Branch if Minus (if N = 1)	REL	18 2B qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBNE <i>rel16</i>	Long Branch if Not Equal (if Z = 0)	REL	18 26 qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBPL <i>rel16</i>	Long Branch if Plus (if N = 0)	REL	18 2A qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBRA <i>rel16</i>	Long Branch Always (if 1=1)	REL	18 20 qq rr	OPPP	OPPP	----	----
LBPN <i>rel16</i>	Long Branch Never (if 1=0)	REL	18 21 qq rr	OPPO	OPPO	----	----
LBVC <i>rel16</i>	Long Branch if Overflow Bit Clear (if V=0)	REL	18 28 qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LBVS <i>rel16</i>	Long Branch if Overflow Bit Set (if V=1)	REL	18 29 qq rr	OPPP/OPPO ¹	OPPP/OPPO ¹	----	----
LDAA <i>oprrB</i> LDAA <i>opraB</i> LDAA <i>oprrBa</i> LDAA <i>opnd_xysp</i> LDAA <i>opnd9_xysp</i> LDAA <i>opnd16_xysp</i> LDAA <i>[D_xysp]</i> LDAA <i>[opnd16_xysp]</i>	(M) → A Load Accumulator A	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	86 11 96 dd 86 hh 11 A6 xb A6 xb ff A6 xb aa ff A6 xb aa ff	P rPF rPO rPF rPO frPP fIFrPF fIFrPF	P rPF rPO rPF rPO frPP fIFrPF fIFrPF	----	AA0-
LDAB <i>oprrB</i> LDAB <i>opraB</i> LDAB <i>oprrBa</i> LDAB <i>opnd_xysp</i> LDAB <i>opnd9_xysp</i> LDAB <i>opnd16_xysp</i> LDAB <i>[D_xysp]</i> LDAB <i>[opnd16_xysp]</i>	(M) → B Load Accumulator B	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	C6 11 D6 dd F6 hh 11 86 xb 86 xb ff 86 xb aa ff 86 xb aa ff	P rPF rPO rPF rPO frPP fIFrPF fIFrPF	P rPF rPO rPF rPO frPP fIFrPF fIFrPF	----	AA0-
LDD <i>oprrB</i> LDD <i>opraB</i> LDD <i>oprrBa</i> LDD <i>opnd_xysp</i> LDD <i>opnd9_xysp</i> LDD <i>opnd16_xysp</i> LDD <i>[D_xysp]</i> LDD <i>[opnd16_xysp]</i>	(M,M+1) → A,B Load Double Accumulator D (A,B)	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	CC jj kk DC dd FC hh 11 8C xb 8C xb ff 8C xb aa ff 8C xb aa ff	PO RPF RPO RPF RPO frPP fIFrPF fIFrPF	OP RPF RPO RPF RPO frPP fIFrPF fIFrPF	----	AA0-

Note 1. OPP/OPPO indicates this instruction takes four cycles to refill the instruction queue if the branch is taken and three cycles if the branch is not taken.

LDS <i>oprrB</i> LDS <i>opraB</i> LDS <i>oprrBa</i> LDS <i>opnd_xysp</i> LDS <i>opnd9_xysp</i> LDS <i>opnd16_xysp</i> LDS <i>[D_xysp]</i> LDS <i>[opnd16_xysp]</i>	(M,M+1) → SP Load Stack Pointer	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	CP jj kk DP dd FP hh 11 8F xb 8F xb ff 8F xb aa ff 8F xb aa ff	OP RPF RPO RPF RPO frPP fIFrPF fIFrPF	OP RPF RPO RPF RPO frPP fIFrPF fIFrPF	----	AA0-
LDX <i>oprrB</i> LDX <i>opraB</i> LDX <i>oprrBa</i> LDX <i>opnd_xysp</i> LDX <i>opnd9_xysp</i> LDX <i>opnd16_xysp</i> LDX <i>[D_xysp]</i> LDX <i>[opnd16_xysp]</i>	(M,M+1) → X Load Index Register X	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	CE jj kk DE dd FE hh 11 8E xb 8E xb ff 8E xb aa ff 8E xb aa ff	OP RPF RPO RPF RPO frPP fIFrPF fIFrPF	OP RPF RPO RPF RPO frPP fIFrPF fIFrPF	----	AA0-

Table A-1. Instruction Set Summary (Sheet 3 of 14)

Source Form	Operation	Addr. Mode	Machine Coding (hex)	Access Detail		SXHI	NZVC
				HCS12	M68HC12		
BLS <i>relB</i>	Branch if Lower or Same (if C + Z = 1) (unsigned)	REL	23 rr	PPP/P ²	PPP/P ²	----	----
BLT <i>relB</i>	Branch if Less Than (if N ⊕ V = 1) (signed)	REL	2D rr	PPP/P ²	PPP/P ²	----	----
BMI <i>relB</i>	Branch if Minus (if N = 1)	REL	2B rr	PPP/P ²	PPP/P ²	----	----
BNE <i>relB</i>	Branch if Not Equal (if Z = 0)	REL	26 rr	PPP/P ²	PPP/P ²	----	----
BPL <i>relB</i>	Branch if Plus (if N = 0)	REL	2A rr	PPP/P ²	PPP/P ²	----	----
BRA <i>relB</i>	Branch Always (if 1 = 1)	REL	20 rr	PPP	PPP	----	----
BACLR <i>opri6a, mskB, relB</i> BACLR <i>opri6a, mskB, relB</i> BACLR <i>opri0_xysp, mskB, relB</i> BACLR <i>opri9_xysp, mskB, relB</i> BACLR <i>opri6_xysp, mskB, relB</i> BACLR <i>opri16_xysp, mskB, relB</i>	Branch if (M) • (mm) = 0 (if All Selected Bit(s) Clear)	DIR EXT IDX IDX1 IDX2	4F dd mm rr 1F hh ll mm rr 0F xb mm rr 0F xb ff mm rr 0F xb aa ff mm rr	rPPP rPPP rPPP rPPP rPPP	rPPP rPPP rPPP rPPP rPPP	----	----
BRN <i>relB</i>	Branch Never (if 1 = 0)	REL	21 rr	P	P	----	----
BRSET <i>opri6a, mskB, relB</i> BRSET <i>opri6a, mskB, relB</i> BRSET <i>opri0_xysp, mskB, relB</i> BRSET <i>opri9_xysp, mskB, relB</i> BRSET <i>opri6_xysp, mskB, relB</i>	Branch if (M) • (mm) = 0 (if All Selected Bit(s) Set)	DIR EXT IDX IDX1 IDX2	4E dd mm rr 1E hh ll mm rr 0E xb mm rr 0E xb ff mm rr 0E xb aa ff mm rr	rPPP rPPP rPPP rPPP rPPP	rPPP rPPP rPPP rPPP rPPP	----	----
BSET <i>opri6, mskB</i> BSET <i>opri6a, mskB</i> BSET <i>opri0_xysp, mskB</i> BSET <i>opri9_xysp, mskB</i> BSET <i>opri6_xysp, mskB</i>	(M) + (mm) → M Set Bit(s) in Memory	DIR EXT IDX IDX1 IDX2	4C dd mm 1C hh ll mm 0C xb mm 0C xb ff mm 0C xb aa ff mm	rPw rPw rPw rPw rPw	rPw rPw rPw rPw rPw	----	AA0-
BSR <i>relB</i>	(SP) - 2 → SP; RTN _L , RTN _H → M _(SP) , M _(SP+1) Subroutine address → PC Branch to Subroutine	REL	07 rr	PPP	PPPS	----	----
BVC <i>relB</i>	Branch if Overflow Bit Clear (if V = 0)	REL	28 rr	PPP/P ²	PPP/P ²	----	----
BVS <i>relB</i>	Branch if Overflow Bit Set (if V = 1)	REL	29 rr	PPP/P ²	PPP/P ²	----	----
CALL <i>opri6a, page</i> CALL <i>opri0_xysp, page</i> CALL <i>opri9_xysp, page</i> CALL <i>opri6_xysp, page</i> CALL <i>[D_xysp]</i> CALL <i>[opri6_xysp]</i>	(SP) - 2 → SP; RTN _L , RTN _H → M _(SP) , M _(SP+1) (SP) - 1 → SP; (PPG) → M _(SP) pg → PPAGE register; Program address → PC Call subroutine in extended memory (Program may be located on another expansion memory page.) Indirect modes get program address and new pg value based on pointer.	EXT IDX IDX1 IDX2	4A hh ll pg 4B xb pg 4B xb ff pg 4B xb aa ff pg	gnSsPPP gnSsPPP gnSsPPP gnSsPPP	gnfSsPPP gnfSsPPP gnfSsPPP gnfSsPPP	----	----
CBA	(A) - (B) Compare 8-Bit Accumulators	INH	18 17	OO	OO	----	AAAA
CLC	0 → C Translates to ANDCC #SFE	IMM	10 FE	P	P	----	---0
CLI	0 → I Translates to ANDCC #SEF (enables I-bit interrupts)	IMM	10 EF	P	P	----	---0
CLR <i>opri6a</i> CLR <i>opri0_xysp</i> CLR <i>opri9_xysp</i> CLR <i>opri6_xysp</i> CLR <i>[D_xysp]</i> CLR <i>[opri6_xysp]</i> CLRA CLRB	0 → M Clear Memory Location 0 → A Clear Accumulator A 0 → B Clear Accumulator B	EXT IDX IDX1 IDX2 [D,IDX] [IDX2] INH INH	79 hh ll 69 xb 69 xb ff 69 xb aa ff 69 xb 69 xb aa ff B7 C7	Pw Pw Pw Pw Pw Pw O O	wP Pw Pw Pw Pw Pw O O	----	0100
CLV	0 → V Translates to ANDCC #FD	IMM	10 FD	P	P	----	--0-
CMPA <i>opri6</i> CMPA <i>opri6a</i> CMPA <i>opri6a</i> CMPA <i>opri0_xysp</i> CMPA <i>opri9_xysp</i> CMPA <i>opri6_xysp</i> CMPA <i>[D_xysp]</i> CMPA <i>[opri6_xysp]</i>	(A) - (M) Compare Accumulator A with Memory	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	B1 11 91 dd B1 hh ll A1 xb A1 xb ff A1 xb aa ff A1 xb A1 xb aa ff	P rP rP rP rP rP rP rP	P rP rP rP rP rP rP rP	----	AAAA

Note 1. PPPP indicates this instruction takes three cycles to refill the instruction queue if the branch is taken and one program fetch cycle if the branch is not taken.

Table A-1. Instruction Set Summary (Sheet 4 of 14)

Source Form	Operation	Addr. Mode	Machine Coding (hex)	Access Detail		SXHI	NZVC
				HCS12	M68HC12		
CMPB #oprB CMPB oprB CMPB opr16a CMPB opr0_xyvsp CMPB opr8_xyvsp CMPB opr16_xyvsp CMPB [D_xyvsp] CMPB [opr16_xyvsp]	(B) - (M) Compare Accumulator B with Memory	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	C1 11 D1 dd F1 hh 11 E1 xb E1 xb ff E1 xb aa ff E1 xb aa ff	P rPF rPO rPF rPO rPF rPF rPF rPF	P rPF rPO rPF rPO rPF rPF rPF	----	AAAA
COM opr16a COM opr0_xyvsp COM opr8_xyvsp COM opr16_xyvsp COM [D_xyvsp] COM [opr16_xyvsp] COMA COMB	(M) → M equivalent to \$FF - (M) → M 1's Complement Memory Location (A) → A Complement Accumulator A (B) → B Complement Accumulator B	EXT IDX IDX1 IDX2 [D,IDX] [IDX2] INH INH	71 hh 11 61 xb 61 xb ff 61 xb aa ff 61 xb aa ff 41 51	rPwO rPw rPwO rPwO rPFwP rPFw rPFw rPFw	rOPw rPw rPwO rPFw rPFw rPFw O O	----	AA01
CPD #opr16 CPD opr16a CPD opr16a CPD opr0_xyvsp CPD opr8_xyvsp CPD opr16_xyvsp CPD [D_xyvsp] CPD [opr16_xyvsp]	(A,B) - (MM+1) Compare D to Memory (16-Bit)	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	8C jj kk 9C dd BC hh 11 AC xb AC xb ff AC xb aa ff AC xb aa ff	PO RPF RPO RPF RPO rRPP rRPF rRPF	CP RPF RDP RPF RPO rRPP rRPF rRPF	----	AAAA
CPS #opr16 CPS opr16a CPS opr16a CPS opr0_xyvsp CPS opr8_xyvsp CPS opr16_xyvsp CPS [D_xyvsp] CPS [opr16_xyvsp]	(SP) - (MM+1) Compare SP to Memory (16-Bit)	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	8F jj kk 9F dd BF hh 11 AF xb AF xb ff AF xb aa ff AF xb aa ff	PO RPF RPO RPF RPO rRPP rRPF rRPF	CP RPF RDP RPF RPO rRPP rRPF rRPF	----	AAAA
CPX #opr16 CPX opr16a CPX opr16a CPX opr0_xyvsp CPX opr8_xyvsp CPX opr16_xyvsp CPX [D_xyvsp] CPX [opr16_xyvsp]	(X) - (MM+1) Compare X to Memory (16-Bit)	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	8E jj kk 9E dd BE hh 11 AE xb AE xb ff AE xb aa ff AE xb aa ff	PO RPF RPO RPF RPO rRPP rRPF rRPF	CP RPF RDP RPF RPO rRPP rRPF rRPF	----	AAAA
CPY #opr16 CPY opr16a CPY opr16a CPY opr0_xyvsp CPY opr8_xyvsp CPY opr16_xyvsp CPY [D_xyvsp] CPY [opr16_xyvsp]	(Y) - (MM+1) Compare Y to Memory (16-Bit)	IMM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	8D jj kk 9D dd BD hh 11 AD xb AD xb ff AD xb aa ff AD xb aa ff	PO RPF RPO RPF RPO rRPP rRPF rRPF	CP RPF RDP RPF RPO rRPP rRPF rRPF	----	AAAA
DAA	Adjust Sum to BCD Decimal Adjust Accumulator A	INH	18 07	OxD	OxD	----	AA?A
DBEQ abdty, rel#	(ctrl) - 1 → ctrl if (ctrl) = 0, then Branch else Continue to next instruction Decrement Counter and Branch if = 0 (ctrl = A, B, D, X, Y, or SP)	REL (9-bit)	04 1b rr	PPP (branch) PP0 (no branch)	PPP	----	----
DBNE abdty, rel#	(ctrl) - 1 → ctrl if (ctrl) not = 0, then Branch; else Continue to next instruction Decrement Counter and Branch if ≠ 0 (ctrl = A, B, D, X, Y, or SP)	REL (9-bit)	04 1b rr	PPP (branch) PP0 (no branch)	PPP	----	----

Source Form	Operation	Address Mode	Machine Coding (Hex)	Access Detail	S X H I N Z V C
STY opr8a STY opr16a STY oprn0_xysspcc STY oprn9_xysspcc STY oprn16_xysspcc STY [D,xysspcc] STY [oprn16,xysspcc]	Store Y (Y _H :Y _L) \rightarrow M:M+1	DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	80 dd 70 hh 11 80 xb 80 xb ff 80 xb 00 ff 80 xb 80 xb 00 ff	rw rwo rw rwo rwp rfrw rfrw	---0000-
SUBA #opr8/ SUBA opr8a SUBA opr16a SUBA oprn0_xysspcc SUBA oprn9_xysspcc SUBA oprn16_xysspcc SUBA [D,xysspcc] SUBA [oprn16,xysspcc]	Subtract from A (A)-(M) \rightarrow A or (A)-imm \rightarrow A	MM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	80 11 90 dd 80 hh 11 A0 xb A0 xb ff A0 xb 00 ff A0 xb A0 xb 00 ff	p rpf rpo rpf rpo frpp frfprf frfprf	---00000
SUBB #opr8/ SUBB opr8a SUBB opr16a SUBB oprn0_xysspcc SUBB oprn9_xysspcc SUBB oprn16_xysspcc SUBB [D,xysspcc] SUBB [oprn16,xysspcc]	Subtract from B (B)-(M) \rightarrow B or (B)-imm \rightarrow B	MM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	80 11 90 dd 80 hh 11 80 xb 80 xb ff 80 xb 00 ff 80 xb 80 xb 00 ff	p rpf rpo rpf rpo frpp frfprf frfprf	---00000
SUBD #opr16/ SUBD opr8a SUBD opr16a SUBD oprn0_xysspcc SUBD oprn9_xysspcc SUBD oprn16_xysspcc SUBD [D,xysspcc] SUBD [oprn16,xysspcc]	Subtract from D (A,B)-(M:M+1) \rightarrow A:B or (A,B)-imm \rightarrow A:B	MM DIR EXT IDX IDX1 IDX2 [D,IDX] [IDX2]	83 jj kk 93 dd 83 hh 11 A3 xb A3 xb ff A3 xb 00 ff A3 xb A3 xb 00 ff	po rpf rpo rpf rpo frpp frfprf frfprf	---00000
SWI	Software interrupt; (SP)-2 \rightarrow SP RTN _i :RTN _L \rightarrow M _{gp} :M _{gp+1} (SP)-2 \rightarrow SP; (Y _H :Y _L) \rightarrow M _{gp} :M _{gp+1} (SP)-2 \rightarrow SP; (X _H :X _L) \rightarrow M _{gp} :M _{gp+1} (SP)-2 \rightarrow SP; (B:A) \rightarrow M _{gp} :M _{gp+1} (SP)-1 \rightarrow SP; (CCR) \rightarrow M _{gp} :1 \rightarrow 1 (SWI vector) \rightarrow PC	INH	3F	vspgpssp*	---1---0
*The CPU also uses vspgpssp for hardware interrupts and unimplemented opcode traps.					
TAB	Transfer A to B; (A) \rightarrow B	INH	18 08	00	---0000-
TAP	Transfer A to CCR; (A) \rightarrow CCR Assembled as TFR A, CCR	INH	87 02	p	00000000
TBA	Transfer B to A; (B) \rightarrow A	INH	18 0F	00	---0000-
TBEQ abcdkysp,rel9	Test and branch if equal to 0 If (counter)=0, then (PC)+2:rel \rightarrow PC	REL (9-bit)	04 1b ff	ppp (branch) rpo (no branch)	---0---0-
TBL oprn0_xysspcc	Table lookup and interpolate, 8-bit (M)+[B] \times (M+1)-(M) \rightarrow A	IDX	18 3D xb	onffff	---00000
TBNE abcdkysp,rel9	Test and branch if not equal to 0 If (counter) \neq 0, then (PC)+2:rel \rightarrow PC	REL (9-bit)	04 1b ff	ppp (branch) rpo (no branch)	---0---0-
TFR abcdkysp,abcdkysp	Transfer from register to register (r1) \rightarrow r2r1 and r2 same size \$00:(r1) \rightarrow r2r1+8-bit; r2=16-bit (r1 _L) \rightarrow r2r1+16-bit; r2=8-bit	INH	87 0b	p	---0---0- or 00000000
TPASame as TFR CCR, A	Transfer CCR to A; (CCR) \rightarrow A	INH	87 20	p	---0---0-

DBNE Decrement and Branch if Not Equal to Zero DBNE

Operation (counter) - 1 ⇒ counter
 If (counter) not = 0, then (PC) + \$0003 + rel ⇒ PC

Subtracts one from the counter register A, B, D, X, Y, or SP. Branches to a relative destination if the counter register does not reach zero. Rel is a 9-bit two's complement offset for branching forward or backward in memory. Branching range is \$100 to \$0FF (-256 to +255) from the address following the last byte of object code in the instruction.

CCR

Effects

S	X	H	I	N	Z	V	C
-	-	-	-	-	-	-	-

Code and CPU Cycles

Source Form	Address Mode	Machine Code (Hex)	CPU Cycles
DBNE <i>abdxysp, rel9</i>	REL (9-bit)	04 1b rr	PPP (branch) PPO (no branch)

Loop Primitive Postbyte (1b) Coding				
Source Form	Postbyte ¹	Object Code	Counter Register	Offset
DBNE A, <i>rel9</i>	0010 X000	04 20 rr	A	Positive
DBNE B, <i>rel9</i>	0010 X001	04 21 rr	B	
DBNE D, <i>rel9</i>	0010 X100	04 24 rr	D	
DBNE X, <i>rel9</i>	0010 X101	04 25 rr	X	
DBNE Y, <i>rel9</i>	0010 X110	04 26 rr	Y	
DBNE SP, <i>rel9</i>	0010 X111	04 27 rr	SP	
DBNE A, <i>rel9</i>	0011 X000	04 30 rr	A	Negative
DBNE B, <i>rel9</i>	0011 X001	04 31 rr	B	
DBNE D, <i>rel9</i>	0011 X100	04 34 rr	D	
DBNE X, <i>rel9</i>	0011 X101	04 35 rr	X	
DBNE Y, <i>rel9</i>	0011 X110	04 36 rr	Y	
DBNE SP, <i>rel9</i>	0011 X111	04 37 rr	SP	

NOTES:

1. Bits 7:6:5 select DBEQ or DBNE; bit 4 is the offset sign bit; bit 3 is not used; bits 2:1:0 select the counter register.

MC9S12 Cycles

- MC9S12 works on **48 MHz clock**
- A processor cycle takes 2 clock cycles – **P** clock is 24 MHz
- Each processor cycle takes **41.7 ns** (1/24 μ s) to execute
- An instruction takes from **1** to **12** processor cycles to execute
- You can determine how many cycles an instruction takes by looking up the CPU cycles for that instruction in the Reference Manual.
 - For example, **LDAA** using the **IMM** addressing mode shows one CPU cycle (of type **P**).
 - **LDAA** using the **EXT** addressing mode shows three CPU cycles (of type **rPO**).
 - Section 6.6 of the S12CPUV2 Reference Manual explains what the HCS12 is doing during each of the different types of CPU cycles.

2000		org \$2000	<i>; Inst</i>	<i>Mode Cycles</i>
2000	C6 0A	ldab #10	<i>; LDAB</i>	<i>(IMM) 1</i>
2002	87	loop: clra	<i>; CLRA</i>	<i>(INH) 1</i>
2003	04 31 FC	dbne b,loop	<i>; DBNE</i>	<i>(REL) 3</i>
2006	3F	swi	<i>; SWI</i>	<i>9</i>

The program executes the **ldab #10** instruction **once** (which takes one cycle). It then goes through the loop **10 times** (which has two instructions, one with one cycle and one with three cycles), and finishes with the swi instruction (which takes 9 cycles).

Total number of cycles:

$$1 + 10 \times (1 + 3) + 9 = 50$$

$$50 \text{ cycles} = 50 \times 41.7 \text{ ns/cycle} = 2.08 \mu\text{s}$$

LDAB

Load B

LDAB

Operation (M) ⇒ B
or
imm ⇒ B

Loads B with either the value in M or an immediate value.

CCR

Effects

S	X	H	I	N	Z	V	C
-	-	-	-	Δ	Δ	0	-

N: Set if MSB of result is set; cleared otherwise

Z: Set if result is \$00; cleared otherwise

V: Cleared

Code and
CPU
Cycles

Source Form	Address Mode	Machine Code (Hex)	CPU Cycles
LDAB #opr8/	IMM	C6 ii	P
LDAB opr8a	DIR	D6 dd	rPF
LDAB opr16a	EXT	F6 hh ll	rPO
LDAB oprx0_xysppc	IDX	B6 xb	rPF
LDAB oprx8_xysppc	IDX1	B6 xb ff	rPO
LDAB oprx16_xysppc	IDX2	B6 xb ee ff	frPP
LDAB [D_xysppc]	[D,IDX]	B6 xb	fifrPF
LDAB [oprx16_xysppc]	[IDX2]	B6 xb ee ff	fIfrPF

Assembler Directives

- In order to write an assembly language program it is necessary to use assembler directives.
- These are not instructions which the HC12 executes but are directives to the assembler program about such things as where to put code and data into memory.
- We will use only a few of these directives. (Note: In the following table, [] means an optional argument.) Here are the ones we will need:

Directive Name	Description	Example
equ	Give a value to a symbol	len: equ 100
org	Set starting value of location counter where code or data will go	org \$1000
dc.b	Allocate and initialize storage for 8-bit variables. Place the bytes in successive memory locations	var: dc.b 2,18 name: dc.b "Jane"
dc.w	Allocate and initialize storage for 16-bit variables. Place the bytes in successive memory locations	var: dc.w \$ABCD
ds.b	Allocate specified number of 8-bit storage places	Table: ds.b 10
ds.w	Allocate specified number of 16-bit storage spaces	table: ds.w 50
dcb.b	Fill memory with a given value The first value is the number of bytes to fill. The second number is the value to put into memory	init_data: dcb.b 100,0

Using labels in assembly programs

A **label** is defined by a name followed by a colon as the first thing on a line. When the label is referred to in the program, it has the numerical value of the location counter when the label was defined.

Here is a code fragment using labels and the assembler directives `dc` and `ds`:

```
org    $2000
table1:dc.b  $23,$17,$f2,$a3,$56
table2: ds.b  5
var:   dc.w  $43af
```

The CodeWarrior assembler produces a listing file (**.lst**). Here is the listing file from the assembler:

```
Freescale HC12-Assembler
(c) Copyright Freescale 1987-2009
Abs.   Rel.   Loc   Obj. code   Source line
----   -
1      1
2      2 a002000 2317 F2A3   table1: dc.b  $23,$17,$f2,$a3,$56
        002004 56
3      3 a002005
4      4 a00200A 43AF     var:   dc.w  $43af
5      5
```

Note that **table1** is a name with the value of \$2000, the value of the location counter defined in the **org** directive. Five bytes of data are defined by the **dc.b** directive, so the location counter is increased from \$2000 to \$2005.

Note that **table2** is a name with the value of \$2005. Five bytes of data are set aside for table2 by the **ds.b 5** directive. The as12 assembler initialized these five bytes of data to all zeros. **var** is a name with the value of \$200a, the first location after table2.

HC12 Instructions

1. Data Transfer and Manipulation Instructions — instructions which move and manipulate data (S12CPUV2 Reference Manual, Sections 5.3, 5.4, and 5.5).

- Load and Store — load copy of memory contents into a register; store copy of register contents into memory.

LDAA \$2000 ; Copy contents of addr \$2000 into A
STD 0,X ; Copy contents of D to addr X and X+1

- Transfer — copy contents of one register to another.

TBA ; Copy B to A
TFR X,Y ; Copy X to Y

- Exchange — exchange contents of two registers.

XGDX ; Exchange contents of D and X
EXG A,B ; Exchange contents of A and B

- Move — copy contents of one memory location to another.

MOVB \$2000,\$20A0 ; Copy byte at \$2000 to \$20A0
MOVW 2,X+,2,Y+ ; Copy two bytes from address held
; in X to address held in Y
; Add 2 to X and Y

2. Arithmetic Instructions — addition, subtraction, multiplication, division (S12CPUV2 Reference Manual, Sections 5.6, 5.8 and 5.12).

ABA ; Add B to A; results in A
SUBD \$20A1 ; Subtract contents of \$20A1 from D
INX ; Increment X by 1
MUL ; Multiply A by B; results in D

3. Logic and Bit Instructions — perform logical operations (S12CPUV2 Reference Manual, Sections 5.9, 5.10, 5.11, 5.13 and 5.14).

- Logic Instructions

ANDA \$2000 ; Logical AND of A with contents of \$2000
EORB 2,X ; Exclusive OR B with contents of address (X+2)

- Clear, Complement and Negate Instructions

NEG -2,X ; Negate (2's comp) contents of address (X-2)
CLRA ; Clear Acc A

- Bit manipulate and test instructions — work with one bit of a register or memory.

BITA # $\$08$; Check to see if Bit 3 of A is set
BSET $\$0002, \#18$; Set bits 3 and 4 of address $\$002$

- Shift and rotate instructions

LSLA ; Logical shift left A
ASR $\$1000$; Arithmetic shift right value at address $\$1000$

4. Compare and test instructions — test contents of a register or memory (to see if zero, negative, etc.), or compare contents of a register to memory (to see if bigger than, etc.) (**S12CPUV2 Reference Manual**, Section 5.9).

TSTA ; (A)-0 -- set flags accordingly
CPX $\#\$8000$; (X) - $\$8000$ -- set flags accordingly

5. Jump and Branch Instructions — Change flow of program (e.g., goto, it-then-else, switch-case) (**S12CPUV2 Reference Manual**, Sections 5.19, 5.20 and 5.21).

JMP L1 ; Start executing code at address label L1
BEQ L2 ; If Z bit set, go to label L2
DBNE X,L3 ; Decrement X; if X not 0 then goto L3
BRCLR $\$1A, \#\$80, L4$; If bit 7 of addr $\$1A$ clear, go to label L4
JSR sub1 ; Jump to subroutine sub1
RTS ; Return from subroutine

6. Interrupt Instructions — Initiate or terminate an interrupt call (**S12CPUV2 Reference Manual**, Section 5.22).

- Interrupt instructions

SWI ; Initiate software interrupt
RTI ; Return from interrupt

7. Index Manipulation Instructions — Put address into X, Y or SP, manipulate X, Y or SP (**S12CPUV2 Reference Manual**, Section 5.23).

ABX ; Add (B) to (X)
LEAX 5,Y ; Put address (Y) + 5 into X

8. Condition Code Instructions — change bits in Condition Code Register (**S12CPUV2 Reference Manual**, Section 5.26).

ANDCC $\#\$f0$; Clear N, Z, C and V bits of CCR
SEV ; Set V bit of CCR

9. Stacking Instructions — push data onto and pull data off of stack (**S12CPUV2 Reference Manual**, Section 5.24).

PSHA ; Push contents of A onto stack
PULX ; Pull two top bytes of stack, put into X

10. Stop and Wait Instructions — put MC9S12 into low power mode (S12CPUV2 Reference Manual, Section 5.27).

STOP ; Put into lowest power mode

WAI ; Put into low power mode until next interrupt

11. Null Instructions

NOP ; No operation

BRN ; Branch never

12. Instructions we won't discuss or use — BCD arithmetic, fuzzy logic, minimum and maximum, multiply-accumulate, table interpolation (**S12CPUV2 Reference Manual**, Sections 5.7, 5.16, 5.17, and 5.18).

Disassembly of an HC12 Program

- It is sometimes useful to be able to convert *HC12 op codes* into *mnemonics*.

For example, consider the hex code:

ADDR DATA

1000 C6 05 CE 20 00 E6 01 18 06 04 35 EE 3F

- To determine the instructions, use Table A-2 of the HCS12 Core Users Guide.
 - If the first byte of the instruction is anything other than **\$18**, use Sheet 1 of Table A.2. From this table, determine the number of bytes of the instruction and the addressing mode. For example, **\$C6** is a two-byte instruction, the mnemonic is **LDAB**, and it uses the **IMM** addressing mode. Thus, the two bytes **C6 05** is the op code for the instruction **LDAB #05**.
 - If the first byte is **\$18**, use Sheet 2 of Table A.2, and do the same thing. For example, **18 06** is a two byte instruction, the mnemonic is **ABA**, and it uses the **INH** addressing mode, so there is no operand. Thus, the two bytes **18 06** is the op code for the instruction **ABA**.
 - Indexed addressing mode is fairly complicated to disassemble. You need to use Table A.3 to determine the operand. For example, the op code **\$E6** indicates **LDAB indexed**, and may use two to four bytes (one to three bytes in addition to the op code). The postbyte **01** indicates that the operand is 0,1, which is **5-bit constant offset**, which takes only one additional byte. All 5-bit constant offset, pre and post increment and decrement, and register offset instructions use one additional byte. All **9-bit constant offset** instructions use two additional bytes, with the second byte holding 8 bits of the 9 bit offset. (**The 9th bit is a direction bit**, which is held in the first postbyte.) All 16-bit constant offset instructions use three postbytes, with the 2nd and 3rd holding the 16-bit unsigned offset.
 - Transfer (**TFR**) and exchange (**EXG**) instructions all have the op code **\$B7**. Use Table A.5 to determine whether it is **TFR** or an **EXG**, and to determine which registers are being used. If the most significant bit of the postbyte is **0**, **the instruction is a transfer instruction**.
 - Loop instructions (Decrement and Branch, Increment and Branch, and Test and Branch) all have the op code **\$04**. To determine which instruction the op code **\$04** implies, and whether the branch is positive (forward) or negative (backward), use Table A.6. For example, in the sequence **04 35 EE**, the 04 indicates a loop

instruction. The 35 indicates it is a **DBNE X** instruction (decrement register X and branch if result is not equal to zero), and the direction is backward (negative). The **EE** indicates a branch of -18 bytes.

- Use up all the bytes for one instruction, then go on to the next instruction.

C6 05	⇒ LDAA #\$05	two-byte LDAA, IMM addressing mode
CE 20 00	⇒ LDX #\$2000	three-byte LDX, IMM addressing mode
E6 01	⇒ LDAB 1,X	two to four-byte LDAB, IDX addressing mode. Operand 01 ⇒ 1,X, a 5b constant offset which uses only one postbyte
18 06	⇒ ABA	two-byte ABA, INH addressing mode
04 35 EE	⇒ DBNE X,(-18)	three-byte loop instruction Postbyte 35 indicates DBNE X, negative
3F	⇒ SWI	one-byte SWI, INH addressing mode

Table A-3. Indexed Addressing Mode Postbyte Encoding (xb)

00	0,X 5b const	10	-16,X 5b const	20	1,+X pre-inc	30	1,X+ post-inc	40	0,Y 5b const	50	-16,Y 5b const	60	1,+Y pre-inc	70	1,Y+ post-inc	80	0,SP 5b const	90	-16,SP 5b const	A0	1,+SP pre-inc	B0	1,SP+ post-inc	C0	0,PC 5b const	D0	-16,PC 5b const	E0	n,X 9b const	F0	n,SP 9b const
01	1,X 5b const	11	-15,X 5b const	21	2,+X pre-inc	31	2,X+ post-inc	41	1,Y 5b const	51	-15,Y 5b const	61	2,+Y pre-inc	71	2,Y+ post-inc	81	1,SP 5b const	91	-15,SP 5b const	A1	2,+SP pre-inc	B1	2,SP+ post-inc	C1	1,PC 5b const	D1	-15,PC 5b const	E1	-n,X 9b const	F1	-n,SP 9b const
02	2,X 5b const	12	-14,X 5b const	22	3,+X pre-inc	32	3,X+ post-inc	42	2,Y 5b const	52	-14,Y 5b const	62	3,+Y pre-inc	72	3,Y+ post-inc	82	2,SP 5b const	92	-14,SP 5b const	A2	3,+SP pre-inc	B2	3,SP+ post-inc	C2	2,PC 5b const	D2	-14,PC 5b const	E2	n,X 16b const	F2	n,SP 16b const
03	3,X 5b const	13	-13,X 5b const	23	4,+X pre-inc	33	4,X+ post-inc	43	3,Y 5b const	53	-13,Y 5b const	63	4,+Y pre-inc	73	4,Y+ post-inc	83	3,SP 5b const	93	-13,SP 5b const	A3	4,+SP pre-inc	B3	4,SP+ post-inc	C3	3,PC 5b const	D3	-13,PC 5b const	E3	[n,X] 16b indir	F3	[n,SP] 16b indir
04	4,X 5b const	14	-12,X 5b const	24	5,+X pre-inc	34	5,X+ post-inc	44	4,Y 5b const	54	-12,Y 5b const	64	5,+Y pre-inc	74	5,Y+ post-inc	84	4,SP 5b const	94	-12,SP 5b const	A4	5,+SP pre-inc	B4	5,SP+ post-inc	C4	4,PC 5b const	D4	-12,PC 5b const	E4	A,X A offset	F4	A,SP A offset
05	5,X 5b const	15	-11,X 5b const	25	6,+X pre-inc	35	6,X+ post-inc	45	5,Y 5b const	55	-11,Y 5b const	65	6,+Y pre-inc	75	6,Y+ post-inc	85	5,SP 5b const	95	-11,SP 5b const	A5	6,+SP pre-inc	B5	6,SP+ post-inc	C5	5,PC 5b const	D5	-11,PC 5b const	E5	B,X B offset	F5	B,SP B offset
06	6,X 5b const	16	-10,X 5b const	26	7,+X pre-inc	36	7,X+ post-inc	46	6,Y 5b const	56	-10,Y 5b const	66	7,+Y pre-inc	76	7,Y+ post-inc	86	6,SP 5b const	96	-10,SP 5b const	A6	7,+SP pre-inc	B6	7,SP+ post-inc	C6	6,PC 5b const	D6	-10,PC 5b const	E6	D,X D offset	F6	D,SP D offset
07	7,X 5b const	17	-9,X 5b const	27	8,+X pre-inc	37	8,X+ post-inc	47	7,Y 5b const	57	-9,Y 5b const	67	8,+Y pre-inc	77	8,Y+ post-inc	87	7,SP 5b const	97	-9,SP 5b const	A7	8,+SP pre-inc	B7	8,SP+ post-inc	C7	7,PC 5b const	D7	-9,PC 5b const	E7	[D,X] D indirect	F7	[D,SP] D indirect
08	8,X 5b const	18	-8,X 5b const	28	8,-X pre-dec	38	8,X- post-dec	48	8,Y 5b const	58	-8,Y 5b const	68	8,-Y pre-dec	78	8,Y- post-dec	88	8,SP 5b const	98	-8,SP 5b const	A8	8,-SP pre-dec	B8	8,SP- post-dec	C8	8,PC 5b const	D8	-8,PC 5b const	E8	n,Y 9b const	F8	n,PC 9b const
09	9,X 5b const	19	-7,X 5b const	29	7,-X pre-dec	39	7,X- post-dec	49	9,Y 5b const	59	-7,Y 5b const	69	7,-Y pre-dec	79	7,Y- post-dec	89	9,SP 5b const	99	-7,SP 5b const	A9	7,-SP pre-dec	B9	7,SP- post-dec	C9	9,PC 5b const	D9	-7,PC 5b const	E9	-n,Y 9b const	F9	-n,PC 9b const
0A	10,X 5b const	1A	-8,X 5b const	2A	6,-X pre-dec	3A	6,X- post-dec	4A	10,Y 5b const	5A	-8,Y 5b const	6A	6,-Y pre-dec	7A	6,Y- post-dec	8A	10,SP 5b const	9A	-8,SP 5b const	AA	6,-SP pre-dec	BA	6,SP- post-dec	CA	10,PC 5b const	DA	-8,PC 5b const	EA	n,Y 16b const	FA	n,PC 16b const
0B	11,X 5b const	1B	-5,X 5b const	2B	5,-X pre-dec	3B	5,X- post-dec	4B	11,Y 5b const	5B	-5,Y 5b const	6B	5,-Y pre-dec	7B	5,Y- post-dec	8B	11,SP 5b const	9B	-5,SP 5b const	AB	5,-SP pre-dec	BB	5,SP- post-dec	CB	11,PC 5b const	DB	-5,PC 5b const	EB	[n,Y] 16b indir	FB	[n,PC] 16b indir
0C	12,X 5b const	1C	-4,X 5b const	2C	4,-X pre-dec	3C	4,X- post-dec	4C	12,Y 5b const	5C	-4,Y 5b const	6C	4,-Y pre-dec	7C	4,Y- post-dec	8C	12,SP 5b const	9C	-4,SP 5b const	AC	4,-SP pre-dec	BC	4,SP- post-dec	CC	12,PC 5b const	DC	-4,PC 5b const	EC	A,Y A offset	FC	A,PC A offset
0D	13,X 5b const	1D	-3,X 5b const	2D	3,-X pre-dec	3D	3,X- post-dec	4D	13,Y 5b const	5D	-3,Y 5b const	6D	3,-Y pre-dec	7D	3,Y- post-dec	8D	13,SP 5b const	9D	-3,SP 5b const	AD	3,-SP pre-dec	BD	3,SP- post-dec	CD	13,PC 5b const	DD	-3,PC 5b const	ED	B,Y B offset	FD	B,PC B offset
0E	14,X 5b const	1E	-2,X 5b const	2E	2,-X pre-dec	3E	2,X- post-dec	4E	14,Y 5b const	5E	-2,Y 5b const	6E	2,-Y pre-dec	7E	2,Y- post-dec	8E	14,SP 5b const	9E	-2,SP 5b const	AE	2,-SP pre-dec	BE	2,SP- post-dec	CE	14,PC 5b const	DE	-2,PC 5b const	EE	D,Y D offset	FE	D,PC D offset
0F	15,X 5b const	1F	-1,X 5b const	2F	1,-X pre-dec	3F	1,X- post-dec	4F	15,Y 5b const	5F	-1,Y 5b const	6F	1,-Y pre-dec	7F	1,Y- post-dec	8F	15,SP 5b const	9F	-1,SP 5b const	AF	1,-SP pre-dec	BF	1,SP- post-dec	CF	15,PC 5b const	DF	-1,PC 5b const	EF	[D,Y] D indirect	FF	[D,PC] D indirect

Key to Table A-3

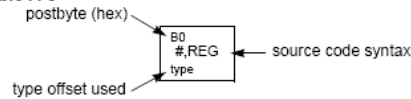


Table A-5. Transfer and Exchange Postbyte Encoding

TRANSFERS									
↓ LS	MS⇒	0	1	2	3	4	5	6	7
0		A ⇒ A	B ⇒ A	CCR ⇒ A	TMP3 _L ⇒ A	B ⇒ A	X _L ⇒ A	Y _L ⇒ A	SP _L ⇒ A
1		A ⇒ B	B ⇒ B	CCR ⇒ B	TMP3 _L ⇒ B	B ⇒ B	X _L ⇒ B	Y _L ⇒ B	SP _L ⇒ B
2		A ⇒ CCR	B ⇒ CCR	CCR ⇒ CCR	TMP3 _L ⇒ CCR	B ⇒ CCR	X _L ⇒ CCR	Y _L ⇒ CCR	SP _L ⇒ CCR
3		sex:A ⇒ TMP2	sex:B ⇒ TMP2	sex:CCR ⇒ TMP2	TMP3 ⇒ TMP2	D ⇒ TMP2	X ⇒ TMP2	Y ⇒ TMP2	SP ⇒ TMP2
4		sex:A ⇒ D SEX A,D	sex:B ⇒ D SEX B,D	sex:CCR ⇒ D SEX CCR,D	TMP3 ⇒ D	D ⇒ D	X ⇒ D	Y ⇒ D	SP ⇒ D
5		sex:A ⇒ X SEX A,X	sex:B ⇒ X SEX B,X	sex:CCR ⇒ X SEX CCR,X	TMP3 ⇒ X	D ⇒ X	X ⇒ X	Y ⇒ X	SP ⇒ X
6		sex:A ⇒ Y SEX A,Y	sex:B ⇒ Y SEX B,Y	sex:CCR ⇒ Y SEX CCR,Y	TMP3 ⇒ Y	D ⇒ Y	X ⇒ Y	Y ⇒ Y	SP ⇒ Y
7		sex:A ⇒ SP SEX A,SP	sex:B ⇒ SP SEX B,SP	sex:CCR ⇒ SP SEX CCR,SP	TMP3 ⇒ SP	D ⇒ SP	X ⇒ SP	Y ⇒ SP	SP ⇒ SP
EXCHANGES									
↓ LS	MS⇒	8	9	A	B	C	D	E	F
0		A ⇔ A	B ⇔ A	CCR ⇔ A	TMP3 _L ⇔ A \$00:A ⇒ TMP3	B ⇔ A A ⇔ B	X _L ⇔ A \$00:A ⇒ X	Y _L ⇔ A \$00:A ⇒ Y	SP _L ⇔ A \$00:A ⇒ SP
1		A ⇔ B	B ⇔ B	CCR ⇔ B	TMP3 _L ⇔ B \$FF:B ⇒ TMP3	B ⇔ B \$FF ⇒ A	X _L ⇔ B \$FF:B ⇒ X	Y _L ⇔ B \$FF:B ⇒ Y	SP _L ⇔ B \$FF:B ⇒ SP
2		A ⇔ CCR	B ⇔ CCR	CCR ⇔ CCR	TMP3 _L ⇔ CCR \$FF:CCR ⇒ TMP3	B ⇔ CCR \$FF:CCR ⇒ D	X _L ⇔ CCR \$FF:CCR ⇒ X	Y _L ⇔ CCR \$FF:CCR ⇒ Y	SP _L ⇔ CCR \$FF:CCR ⇒ SP
3		\$00:A ⇒ TMP2 TMP2 _L ⇒ A	\$00:B ⇒ TMP2 TMP2 _L ⇒ B	\$00:CCR ⇒ TMP2 TMP2 _L ⇒ CCR	TMP3 ⇔ TMP2	D ⇔ TMP2	X ⇔ TMP2	Y ⇔ TMP2	SP ⇔ TMP2
4		\$00:A ⇒ D	\$00:B ⇒ D	\$00:CCR ⇒ D B ⇔ CCR	TMP3 ⇔ D	D ⇔ D	X ⇔ D	Y ⇔ D	SP ⇔ D
5		\$00:A ⇒ X X _L ⇒ A	\$00:B ⇒ X X _L ⇒ B	\$00:CCR ⇒ X X _L ⇒ CCR	TMP3 ⇔ X	D ⇔ X	X ⇔ X	Y ⇔ X	SP ⇔ X
6		\$00:A ⇒ Y Y _L ⇒ A	\$00:B ⇒ Y Y _L ⇒ B	\$00:CCR ⇒ Y Y _L ⇒ CCR	TMP3 ⇔ Y	D ⇔ Y	X ⇔ Y	Y ⇔ Y	SP ⇔ Y
7		\$00:A ⇒ SP SP _L ⇒ A	\$00:B ⇒ SP SP _L ⇒ B	\$00:CCR ⇒ SP SP _L ⇒ CCR	TMP3 ⇔ SP	D ⇔ SP	X ⇔ SP	Y ⇔ SP	SP ⇔ SP

TMP2 and TMP3 registers are for factory use only.

Table A-6. Loop Primitive Postbyte Encoding (1b)

00	A	DBEQ (+)	10	A	DBEQ (-)	20	A	DBNE (+)	30	A	DBNE (-)	40	A	TBEQ (+)	50	A	TBEQ (-)	60	A	TBNE (+)	70	A	TBNE (-)	80	A	IBEQ (+)	90	A	IBEQ (-)	A0	A	IBNE (+)	B0	A	IBNE (-)
01	B	DBEQ (+)	11	B	DBEQ (-)	21	B	DBNE (+)	31	B	DBNE (-)	41	B	TBEQ (+)	51	B	TBEQ (-)	61	B	TBNE (+)	71	B	TBNE (-)	81	B	IBEQ (+)	91	B	IBEQ (-)	A1	B	IBNE (+)	B1	B	IBNE (-)
02		—	12		—	22		—	32		—	42		—	52		—	62		—	72		—	82		—	92		—	A2		—	B2		—
03		—	13		—	23		—	33		—	43		—	53		—	63		—	73		—	83		—	93		—	A3		—	B3		—
04	D	DBEQ (+)	14	D	DBEQ (-)	24	D	DBNE (+)	34	D	DBNE (-)	44	D	TBEQ (+)	54	D	TBEQ (-)	64	D	TBNE (+)	74	D	TBNE (-)	84	D	IBEQ (+)	94	D	IBEQ (-)	A4	D	IBNE (+)	B4	D	IBNE (-)
05	X	DBEQ (+)	15	X	DBEQ (-)	25	X	DBNE (+)	35	X	DBNE (-)	45	X	TBEQ (+)	55	X	TBEQ (-)	65	X	TBNE (+)	75	X	TBNE (-)	85	X	IBEQ (+)	95	X	IBEQ (-)	A5	X	IBNE (+)	B5	X	IBNE (-)
06	Y	DBEQ (+)	16	Y	DBEQ (-)	26	Y	DBNE (+)	36	Y	DBNE (-)	46	Y	TBEQ (+)	56	Y	TBEQ (-)	66	Y	TBNE (+)	76	Y	TBNE (-)	86	Y	IBEQ (+)	96	Y	IBEQ (-)	A6	Y	IBNE (+)	B6	Y	IBNE (-)
07	SP	DBEQ (+)	17	SP	DBEQ (-)	27	SP	DBNE (+)	37	SP	DBNE (-)	47	SP	TBEQ (+)	57	SP	TBEQ (-)	67	SP	TBNE (+)	77	SP	TBNE (-)	87	SP	IBEQ (+)	97	SP	IBEQ (-)	A7	SP	IBNE (+)	B7	SP	IBNE (-)

Key to Table A-6

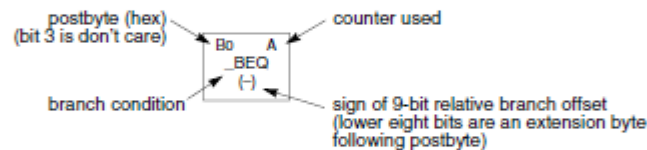


Table A-7. Branch/Complementary Branch

Branch				Complementary Branch			
Test	Mnemonic	Opcode	Boolean	Test	Mnemonic	Opcode	Comment
r>m	BGT	2E	Z + (N ⊕ V) = 0	r≤m	BLE	2F	Signed
r≥m	BGE	2C	N ⊕ V = 0	r<m	BLT	2D	Signed
r=m	BEQ	27	Z = 1	r≠m	BNE	26	Signed
r≤m	BLE	2F	Z + (N ⊕ V) = 1	r>m	BGT	2E	Signed
r<m	BLT	2D	N ⊕ V = 1	r≥m	BGE	2C	Signed
r>m	BHI	22	C + Z = 0	r≤m	BLS	23	Unsigned
r≥m	BHS/BCC	24	C = 0	r<m	BLO/BCS	25	Unsigned
r=m	BEQ	27	Z = 1	r≠m	BNE	26	Unsigned
r≤m	BLS	23	C + Z = 1	r>m	BHI	22	Unsigned
r<m	BLO/BCS	25	C = 1	r≥m	BHS/BCC	24	Unsigned
Carry	BCS	25	C = 1	No Carry	BCC	24	Simple
Negative	BMI	2B	N = 1	Plus	BPL	2A	Simple
Overflow	BVS	29	V = 1	No Overflow	BVC	28	Simple
r=0	BEQ	27	Z = 1	r≠0	BNE	26	Simple
Always	BRA	20	—	Never	BRN	21	Unconditional

For 16-bit offset long branches precede opcode with a \$18 page prebyte.