

THE STACK AND THE STACK POINTER

- Sometimes it is useful to have a region of memory for temporary storage, which does not have to be allocated as named variables.
- When we use subroutines and interrupts it will be essential to have such a storage region.
- Such a region is called a *Stack*.
- The *Stack Pointer* (SP) register is used to indicate the location of the last item put onto the stack.
- When you put something onto the stack (push onto the stack), the SP is decremented *before* the item is placed on the stack.
- When you take something off of the stack (pull from the stack), the SP is incremented *after* the item is pulled from the stack.
- Before you can use a stack you have to initialize the Stack Pointer to point to one value higher than the highest memory location in the stack.
- For the HC12 use a block of memory from about \$3B00 to \$3BFF for the stack.
- For this region of memory, initialize the stack pointer to \$3C00.
- Use the LDS (Load Stack Pointer) instruction to initialize the stack point.
- The LDS instruction is usually the first instruction of a program which uses the stack.
- The stack pointer is initialized only one time in the program.
- For microcontrollers such as the HC12, it is up to the programmer to know how much stack his/her program will need, and to make sure enough space is allocated for the stack. If not enough space is allocated the stack can overwrite data and/or code, which will cause the program to malfunction or crash.

The stack is an array of memory dedicated to temporary storage

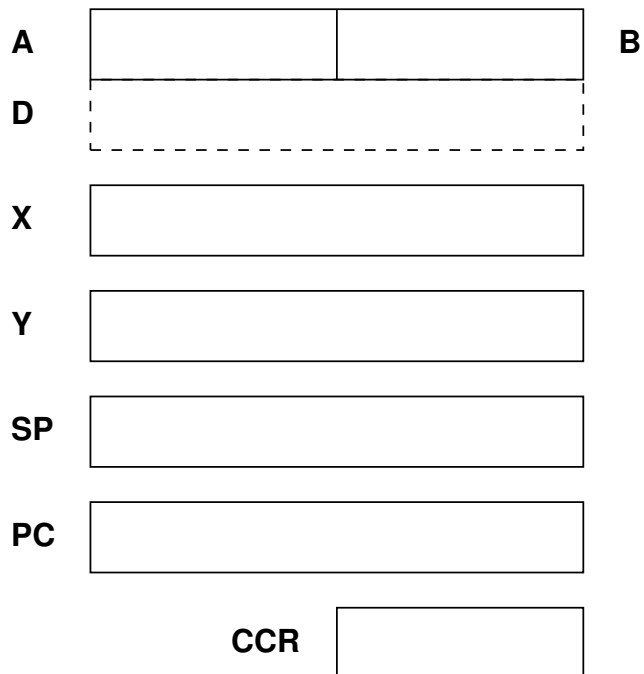
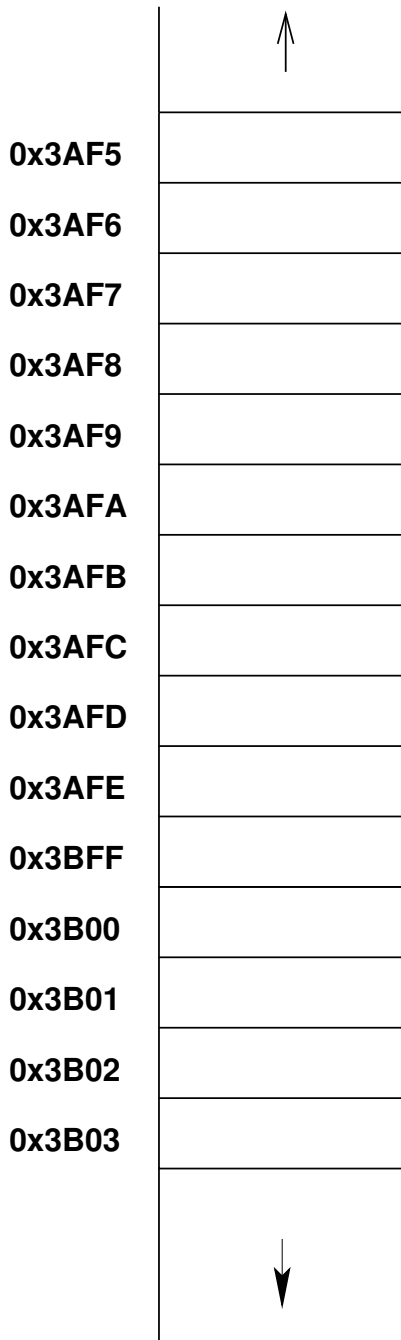
**SP points to location last item
placed in block**

SP decreases when you put item on stack

SP increases when you pull item from stack

For HC12 EVBU, use 0x3C00 as initial SP:

```
STACK:  EQU  $3C00
        LDS  #STACK
```



An example of some code which uses the stack

Stack Pointer:

Initialize ONCE before first use (LDS #STACK)

Points to last used storage location

Decreases when you put something on stack

Increases when you take something off stack



STACK: equ \$3C00

CODE: section .text
org 0x1000

```
lds    #STACK
ldaa   #$2e
ldx    #$1254
psha
pshx
clra
ldx    #$ffff
```

CODE THAT USES A & X

```
pulx
pula
```

A

X

SP

PSHA

Push A onto Stack

PSHA

Operation

$(SP) - \$0001 \Rightarrow SP$
 $(A) \Rightarrow M_{SP}$

Decrements SP by one and loads the value in A into the address to which SP points.

Push instructions are commonly used to save the contents of one or more CPU registers at the start of a subroutine. Complementary pull instructions can be used to restore the saved CPU registers just before returning from the subroutine.

CCR

Effects

S	X	H	I	N	Z	V	C
-	-	-	-	-	-	-	-

Code and CPU Cycles

Source Form	Address Mode	Machine Code (Hex)	CPU Cycles
PSHA	INH	36	0s

Subroutines

- A subroutine is a section of code which performs a specific task, usually a task which needs to be executed by different parts of a program.
- Example:
 - Math functions, such as *square root*
- Because a subroutine can be called from different places in a program, you cannot get out of a subroutine with an instruction such as

`jmp label`

because you would need to jump to different places depending upon which section of code called the subroutine.

- When you want to call the subroutine your code has to save the address where the subroutine should return to. It does this by saving the *return address* on the stack.
 - This is done automatically for you when you get to the subroutine by using the JSR (Jump to Subroutine) or BSR (Branch to Subroutine) instruction. This instruction pushes the address of the instruction following the JSR (BSR) instruction on the stack.
- After the subroutine is done executing its code it needs to return to the address saved on the stack.
 - This is done automatically for you when you return from the subroutine by using the RTS (Return from Subroutine) instruction. This instruction pulls the return address off of the stack and loads it into the program counter, so the program resumes execution of the program with the instruction following that which called the subroutine.

The subroutine will probably need to use some HC12 registers to do its work. However, the calling code may be using its registers for some reason — the calling code may not work correctly if the subroutine changes the values of the HC12 registers.

- To avoid this problem, the subroutine should save the HC12 registers before it uses them, and restore the HC12 registers after it is done with them.

BSR

Branch to Subroutine

BSR

Operation

$(SP) - \$0002 \Rightarrow SP$
 $RTN_H:RTN_L \Rightarrow M_{SP}:M_{SP+1}$
 $(PC) + \$0002 + rel \Rightarrow PC$

Sets up conditions to return to normal program flow, then transfers control to a subroutine. Uses the address of the instruction after the BSR as a return address.

Decrements the SP by two, to allow the two bytes of the return address to be stacked.

Stacks the return address (the SP points to the high byte of the return address).

Branches to a location determined by the branch offset.

Subroutines are normally terminated with an RTS instruction, which restores the return address from the stack.

CCR

Effects

S	X	H	I	N	Z	V	C
–	–	–	–	–	–	–	–

Code and CPU Cycles

Source Form	Address Mode	Machine Code (Hex)	CPU Cycles
BSR <i>rel8</i>	REL	07 <i>rr</i>	SPPP

RTS

Return from Subroutine

RTS

Operation $(M_{SP}): (M_{SP+1}) \Rightarrow PC_H:PC_L$
 $(SP) + \$0002 \Rightarrow SP$

Restores the value of PC from the stack and increments SP by two. Program execution continues at the address restored from the stack.

CCR								
Effects	S	X	H	I	N	Z	V	C
	-	-	-	-	-	-	-	-

Code and CPU Cycles	Source Form	Address Mode	Machine Code (Hex)	CPU Cycles
	RTS	INH	3D	UfPPP

Example of a subroutine to delay for a certain amount of time

; Subroutine to wait for 100 ms


```
delay:  ldaa    #250
loop2:  ldx     #800
loop1:  dex
        bne     loop1
        deca
        bne     loop2
        rts
```

- Problem: The subroutine changes the values of registers A and X
- To solve, save the values of A and X on the stack before using them, and restore them before returning.

; Subroutine to wait for 100 ms

```
delay:  psha                ; Save regs used by sub on stack
        pshx
        ldaa    #250
loop2:  ldx     #800
loop1:  dex
        bne     loop1
        deca
        bne     loop2
        pulx                ; Restore regs in opposite
        pula                ; order
        rts
```

```
; Program to make a binary counter on LEDs
;
; The program uses a subroutine to insert a delay
; between counts

prog:    equ        $1000
STACK:   equ        $3C00        ;Stack ends of $3BFF
PORTA:   equ        $0000
PORTB:   equ        $0001
DDRA:    equ        $0002
DDRB:    equ        $0003

        org        prog

        lds        #STACK        ; initialize stack pointer
        ldaa       #$ff          ; put all ones into DDRA
        staa       DDRA          ; to make PORTA output
        clr        PORTA         ; put $00 into PORTA
loop:    jsr        delay         ; wait a bit
        inc        PORTA         ; add one to PORTA
        bra        loop          ; repeat forever

; Subroutine to wait for 100 ms

delay:   psha
        pshx
        ldaa       #250
loop2:   ldx        #800
loop1:   dex
        bne        loop1
        deca
        bne        loop2
        pulx
        pula
        rts
```


- Wait until bit 7 of address \$00C4 is set.
- Write the value in ACCA to address \$00C7.

```
; This routine waits until the HC12 serial
; port is ready, then sends a byte of data
; to the HC12 serial port
```

```
putchar:      brclr    $00CC,$$80,putchar
              staa     $00CF
              rts
```

- Program to send the word `hello` to the HC12 serial port

```
; Program fragment to write the word "hello" to the
; HC12 serial port
```

```
              ldx      $str
loop:         ldaa     1,x+    ; get next char
              beq      done    ; char == 0 => no more
              jsr      putchar
              bra      loop
              swi

str:          dc.b     "hello"
              fc.b     $0A,$0D,0    ; CR LF
```

Here is the complete program to write a line to the screen:

```
prog:    equ    $1000
data:    equ    $2000
stack:   equ    $3c00

        org    prog
        lds    #stack
        ldx    #str
loop:    ldaa    1,x+ ; get next char
        beq    done ; char == 0 => no more
        jsr    putchar
        bra    loop
done:    swi

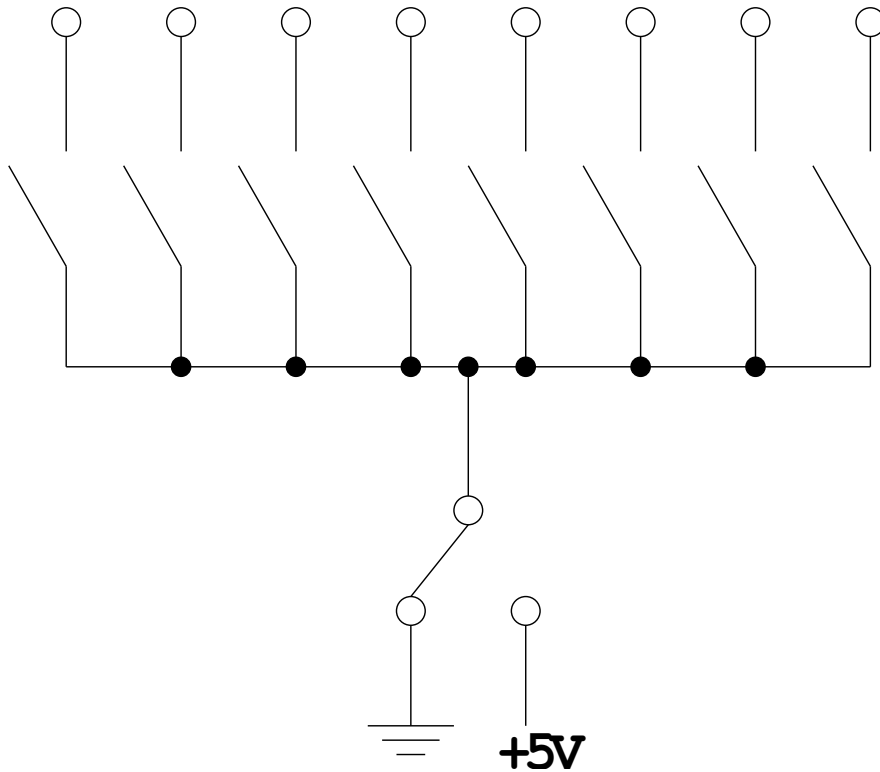
putchar: brclr    $00CC,$80,putchar
        staa    $00CF
        rts

        org    data
str:     fcc     "hello"
        dc.b    $0a,$0d,0 ; CR LF
```

Using DIP switches to get data into the HC12

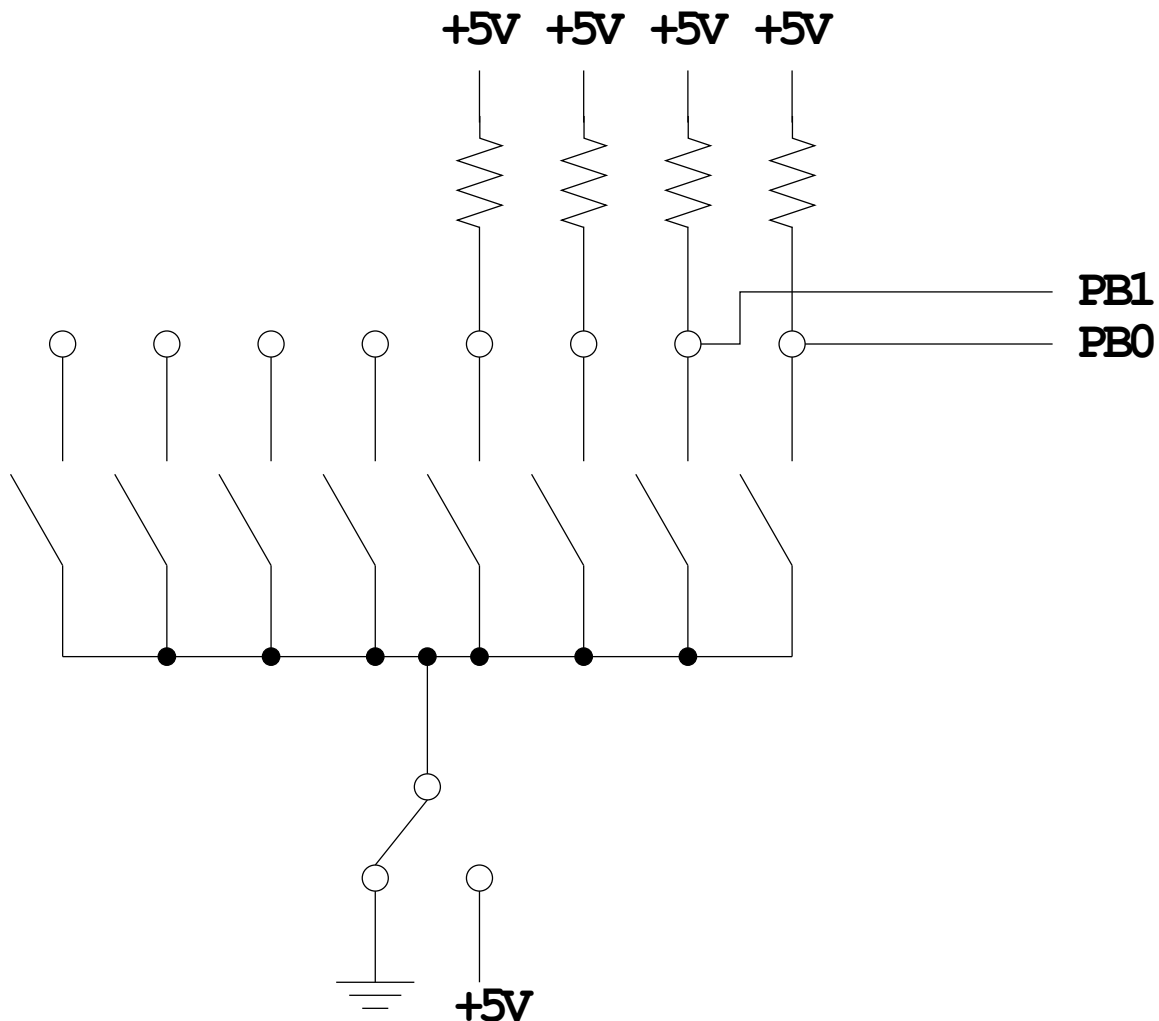
- DIP switches make or break a connection (usually to ground)

DIP Switches on Breadboard



- To use DIP switches, connect one end of each switch to a resistor
- Connect the other end of the resistor to +5 V
- Connect the junction of the DIP switch and the resistor to an input port on the HC12

Using DIP Switches



- When the switch is open, the input port sees a logic 1 (+5 V)
- When the switch is closed, the input sees a logic 0 (0 V)

Looking at the state of a few input pins

- Want to look for a particular pattern on 4 input pins
 - For example want to do something if pattern on PB3-PB0 is 0110
- Don't know or care what are on the other 4 pins (PB7-PB4)
- Here is the wrong way to do it:

```
ldaa    PORTB
cmpa    #b0110
beq     task
```

- If PB7-PB4 are anything other than 0000, you will not execute the task.
- You need to mask out the Don't Care bits **before** checking for the pattern on the bits you are interested in

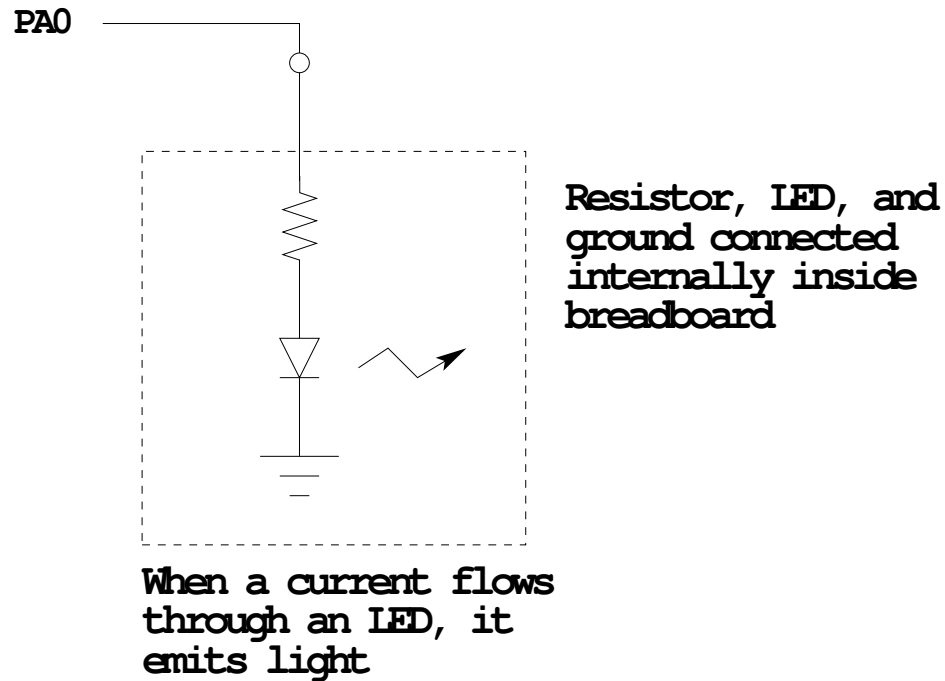
```
ldaa    PORTB
anda    #b00001111
cmpa    #b00000110
beq     task
```

- Now, whatever pattern appears on PB7-4 is ignored

Using an HC12 output port to control an LED

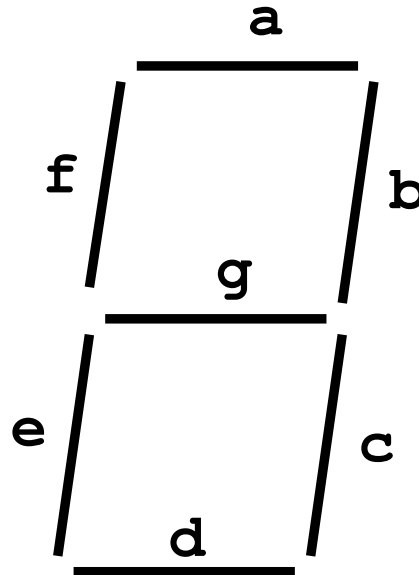
- Connect an output port from the HC12 to an LED.

Using an output port to control an LED



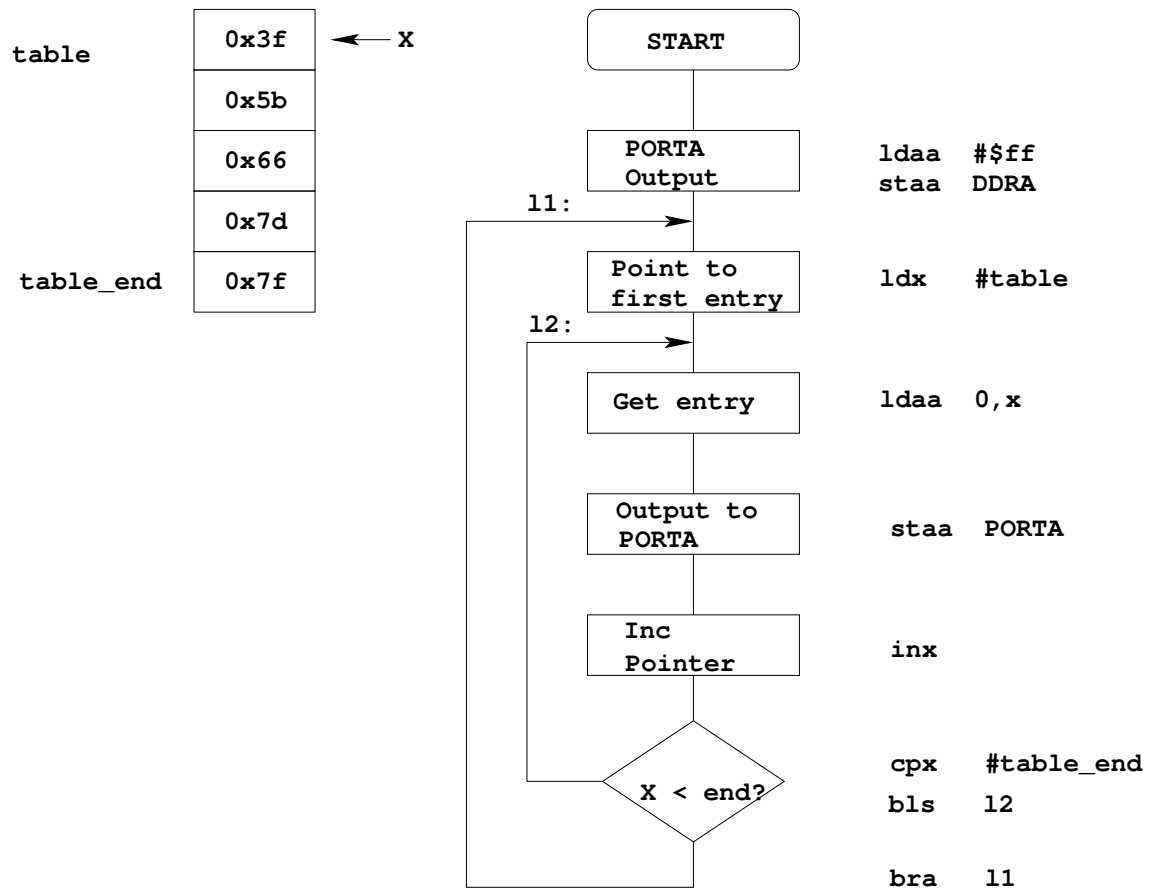
Making a pattern on a seven-segment LED

- Want to generate a particular pattern on a seven-segment LED:



- Determine a number (hex or binary) which will generate each element of the pattern
 - For example, to display a 0, turn on segments a, b, c, d, e and f, or bits 0, 1, 2, 3, 4 and 5 of PTH. The binary pattern is 00111111, or \$3f.
 - To display 0 2 4 6 8, the hex numbers are \$3f, \$5b, \$66, \$7d, \$7f.
- Put the numbers in a table
- Go through the table one by one to display the pattern
- When you get to the last element, repeat the loop

Flowchart to display a pattern of lights on a set of LEDs



; Program using subroutine to make a time delay

```

prog:      equ      $1000
data:      equ      $2000
stack:     equ      $3C00
PTH:       equ      $0260
DDRH:      equ      $0262

          org      prog

          lds      #stack      ; initialize stack pointer
          ldaa     #$ff        ; Make PTH output
          staa     DDRH        ; 0xFF -> DDRH
11:        ldx     #table      ; Start pointer at table
12:        ldaa     1,x+        ; Get value; point to next
          staa     PTH         ; Update LEDs
          jsr      delay       ; Wait a bit
          cpx      #table_end   ; More to do?
          bls      12          ; Yes, keep going through table
          bra      11          ; At end; reset pointer

delay:     psha
          pshx
          ldaa     #250
loop2:     ldx     #8000
loop1:     dex
          bne      loop1
          deca
          bne      loop2
          pulx
          pula
          rts

          org      data
table:     dc.b     $3f
          dc.b     $5b
          dc.b     $66
          dc.b     $7d
table_end: dc.b     $7F

```